Spatially Distributed Queues

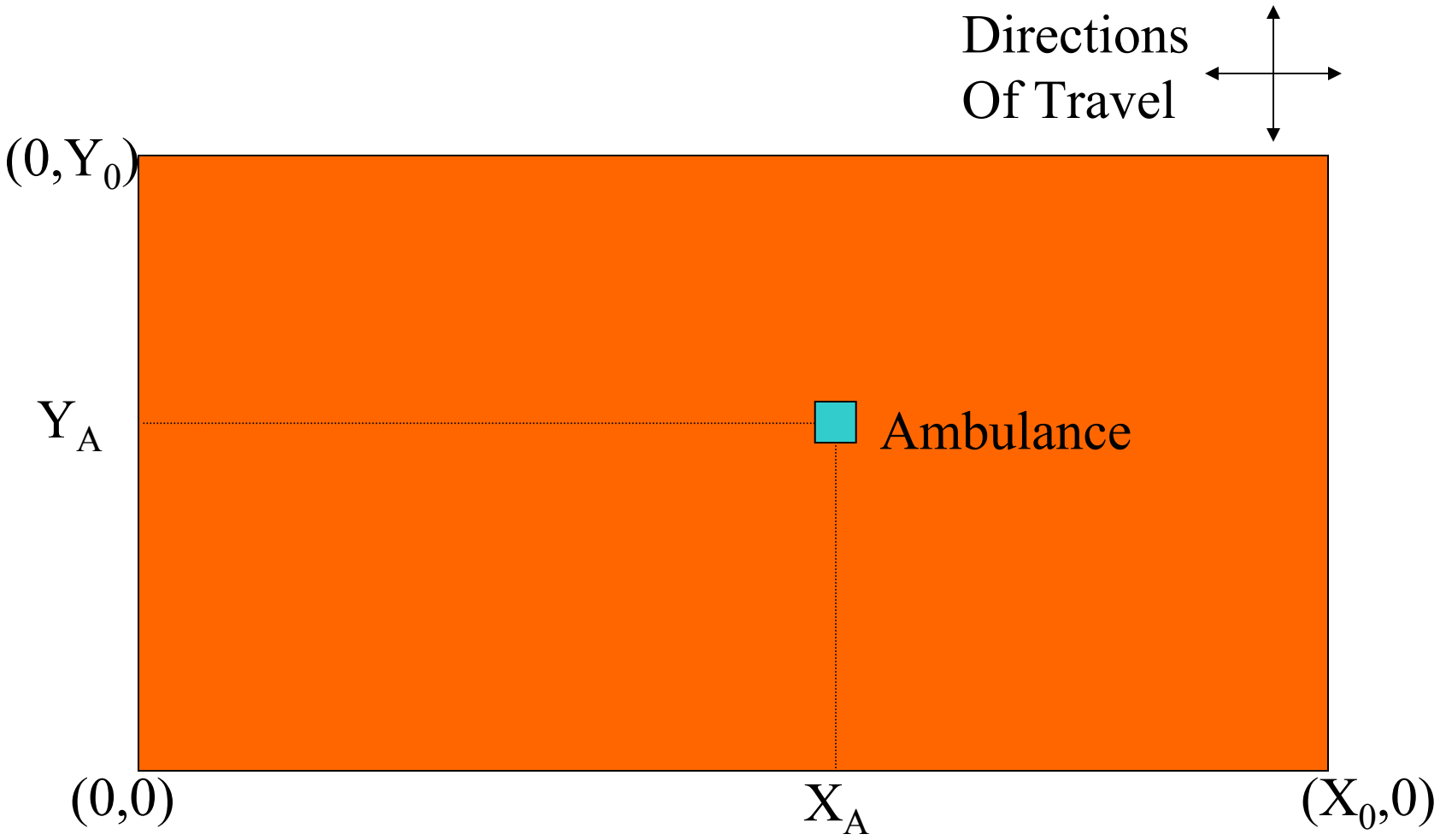
M/G/1

2 Servers

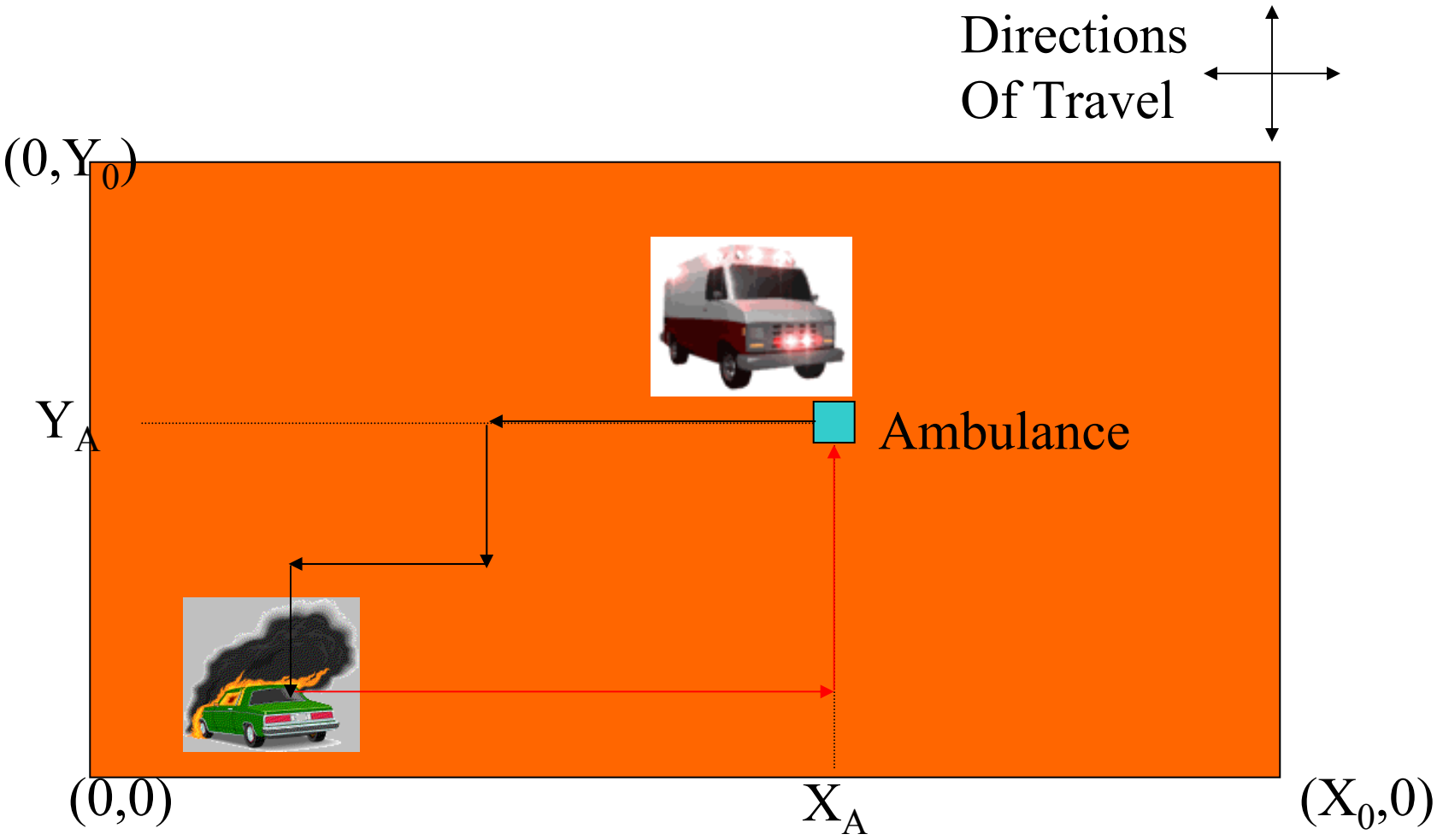
N servers

Approximations

M/G/1



M/G/1



M/G/1

- Ambulance always returns home with each service; standard M/G/1 applies
- But suppose we have an emergency repair vehicle that travels directly from one customer to the next?……

M/G/1with different 1st service time

$\overline{S}_1, \sigma^2_{S_1}$ = expected value and variance,respectively,

of the 1st service time in a busy period

$\overline{S}_2, \sigma^2_{S_2}$ = expected value and variance,respectively,

of the 2nd & all succeeding service times in a busy period

$\lambda \overline{S}_2 < 1$

$\rho = 1 - P_0$ = fraction of time server is busy

M/G/1with different 1st service time

$$\rho = \frac{\lambda \bar{S}_1}{1 - \lambda(\bar{S}_2 - \bar{S}_1)}$$

$$L = \rho + \frac{\lambda^2}{1 - \lambda(\bar{S}_2 - \bar{S}_1)} \left[\frac{\sigma_{S_1}^2 + \bar{S}_1^2 + \lambda\{\bar{S}_1(\sigma_{S_2}^2 + \bar{S}_2^2) - \bar{S}_2(\sigma_{S_1}^2 + \bar{S}_1^2)}{2(1 - \lambda \bar{S}_2)} \right]$$

M/G/1with different 1st service time

Little' s Law : Buy one, get three others for free!

$$L = \lambda W$$

$$L_q = \lambda W_q$$

See the book, Eqs. (5.0) - (5.5)

M/G/1with different 1st service time

- Does this new more general M/G/1 model apply exactly to the ambulance problem?

- Why or why not?

Two-Server “Hypercube” Queueing Model

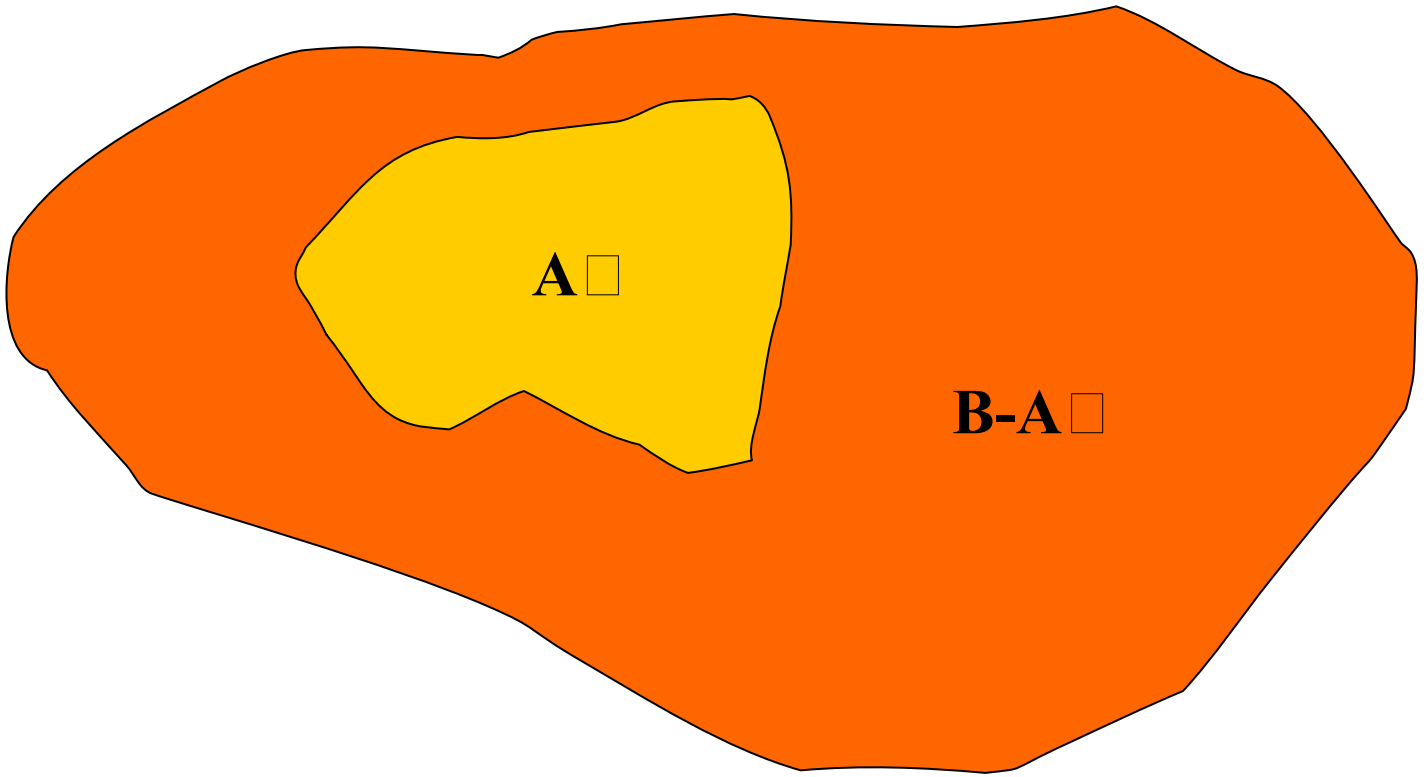
- Distinguishable servers
- Different workloads (due to geography)
- Can appear with or without queueing
  - With -- usually FCFS
  - Without -- usually means a backup contract service is in place


119
瀬戸市
消防本部


119
瀬戸市
消防本部

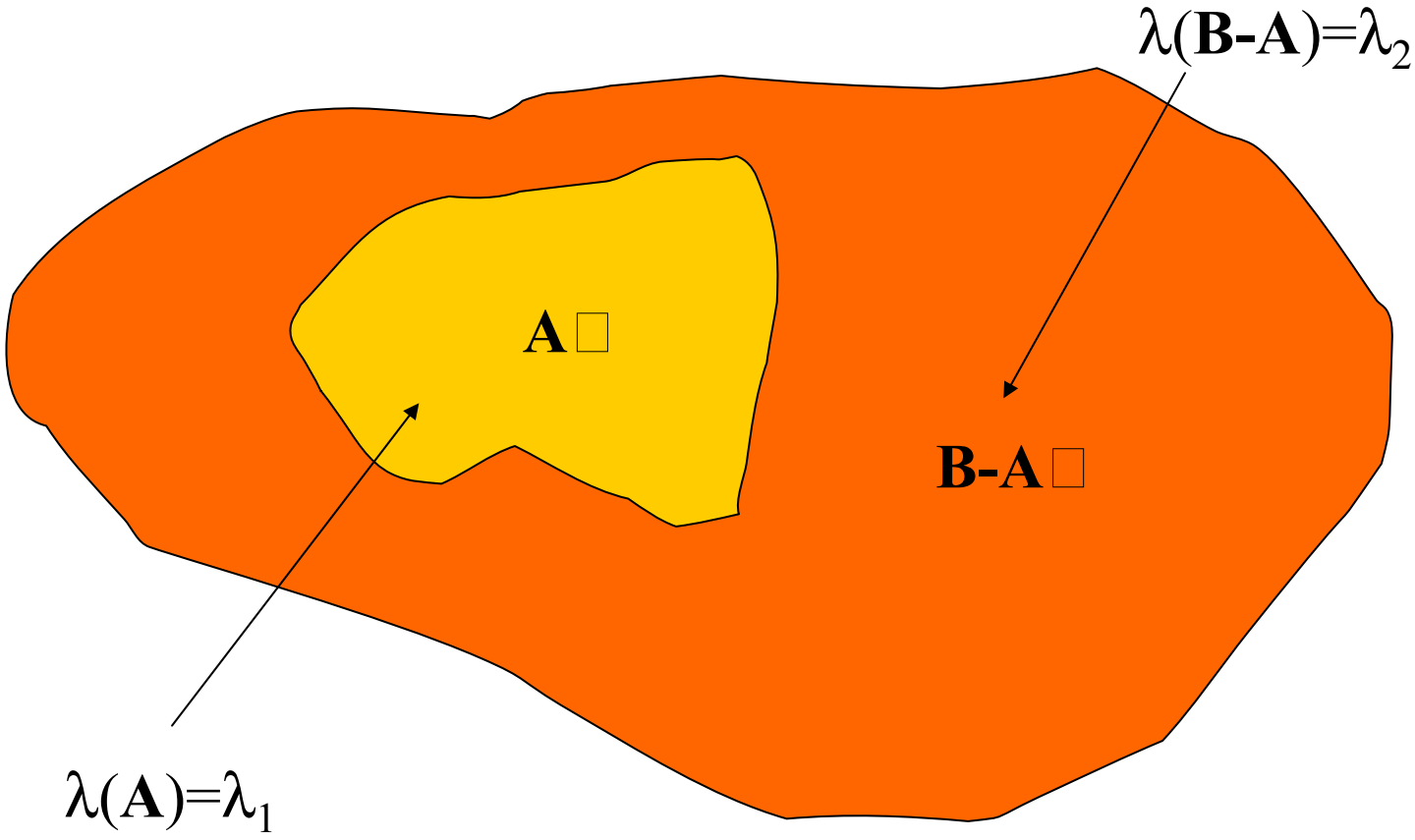

FIRE DEPARTMENT
SEXTO


FIRE DEPARTMENT
SEXTO



B = “Service Region”

Poisson Arrivals from any sub-region **A**

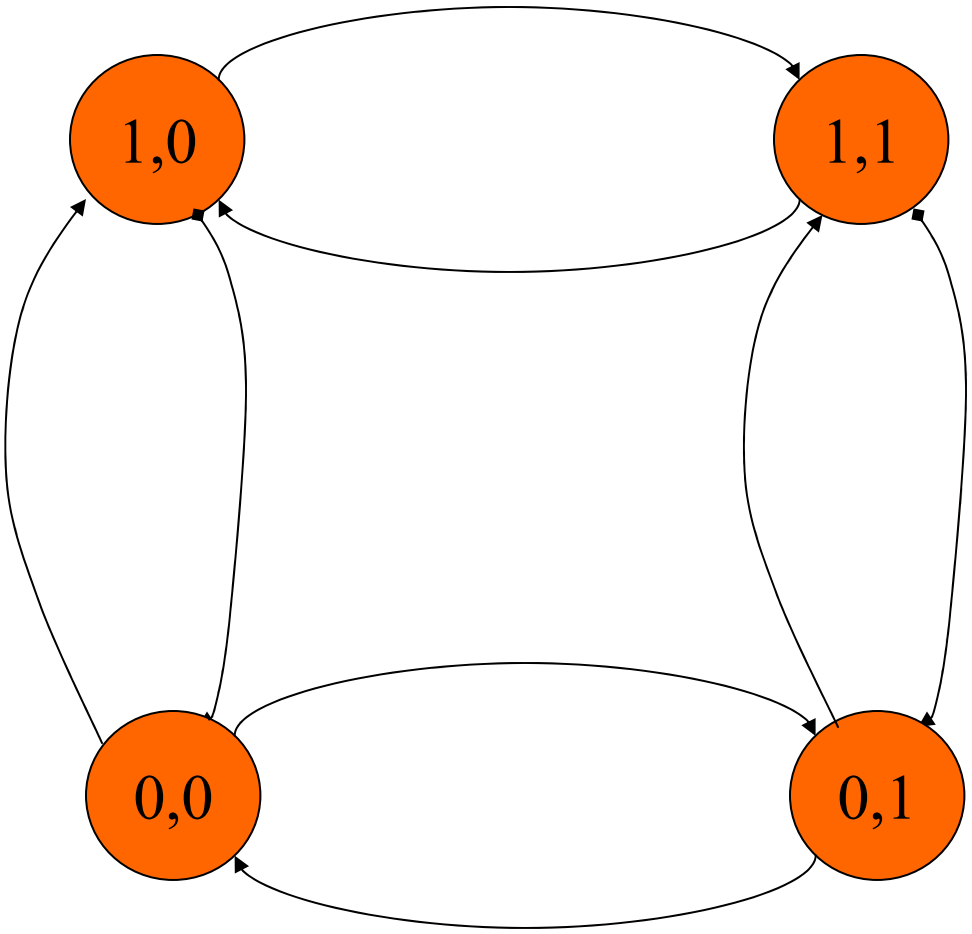


B = "Service Region"

$\lambda = \lambda_1 + \lambda_2$

Service Discipline

- 1st Dispatch Preference to ‘primary server’
- Otherwise, assign customer to other server, if available
- Otherwise, job is ‘lost” (What happens in practice?)


1,0
1,1
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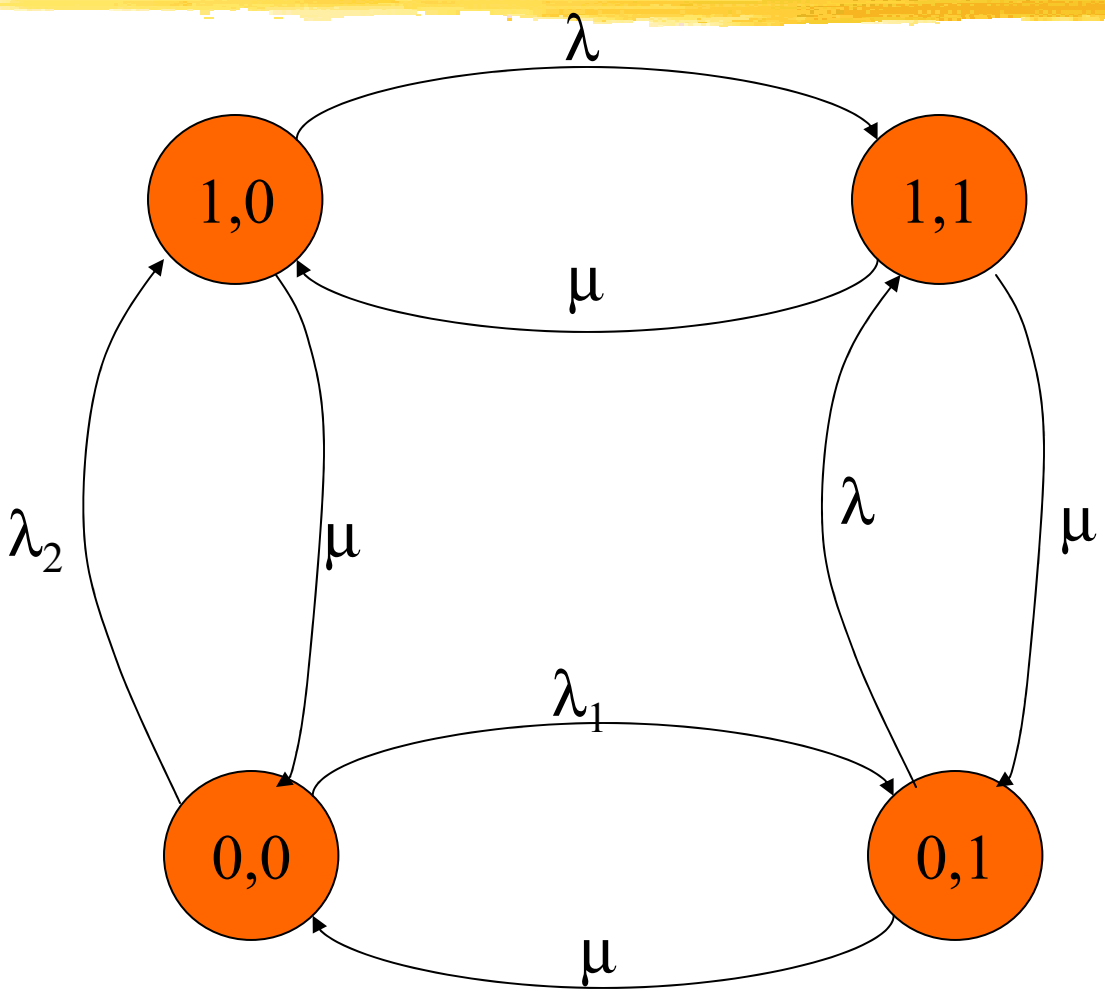
λ
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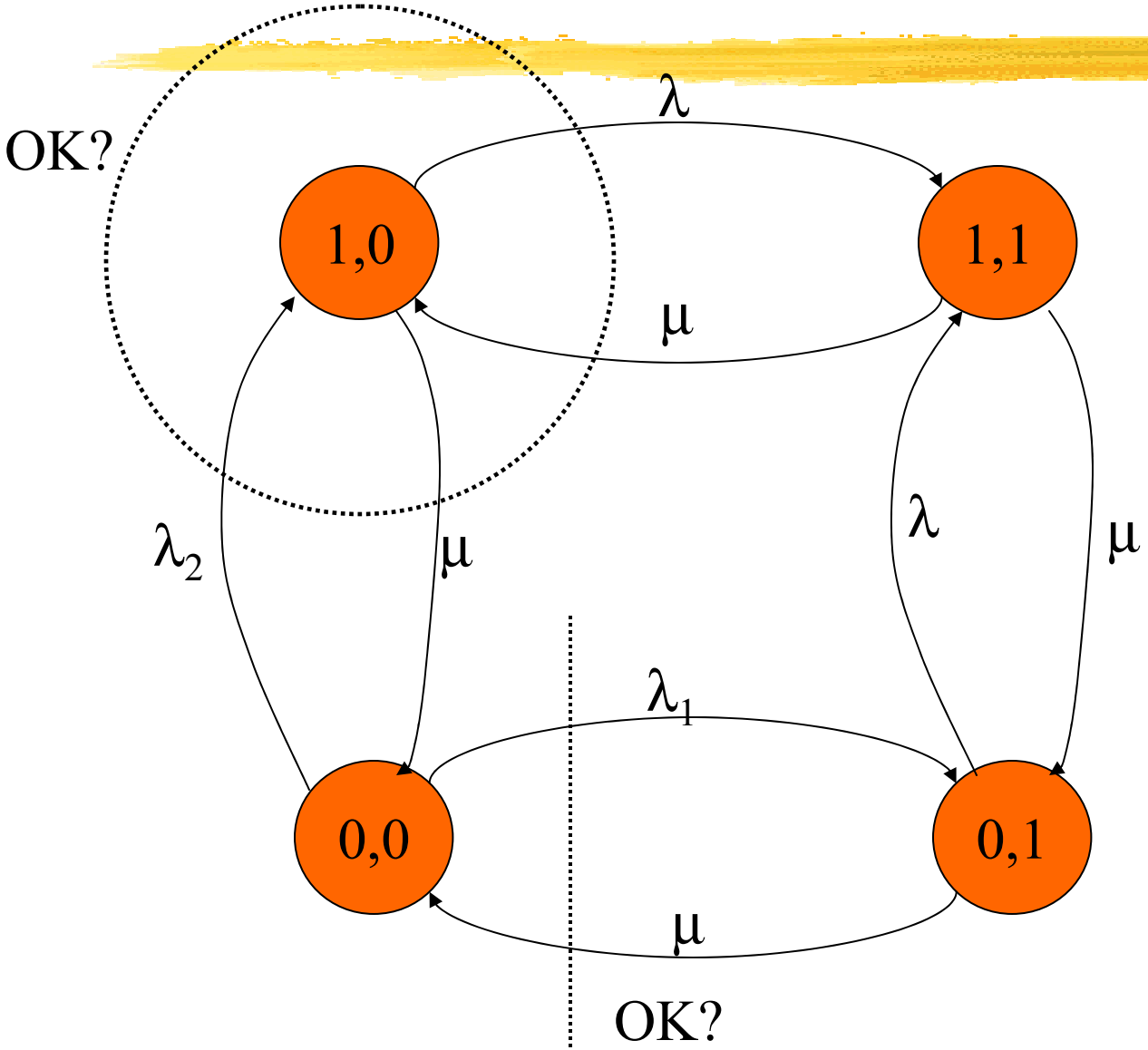
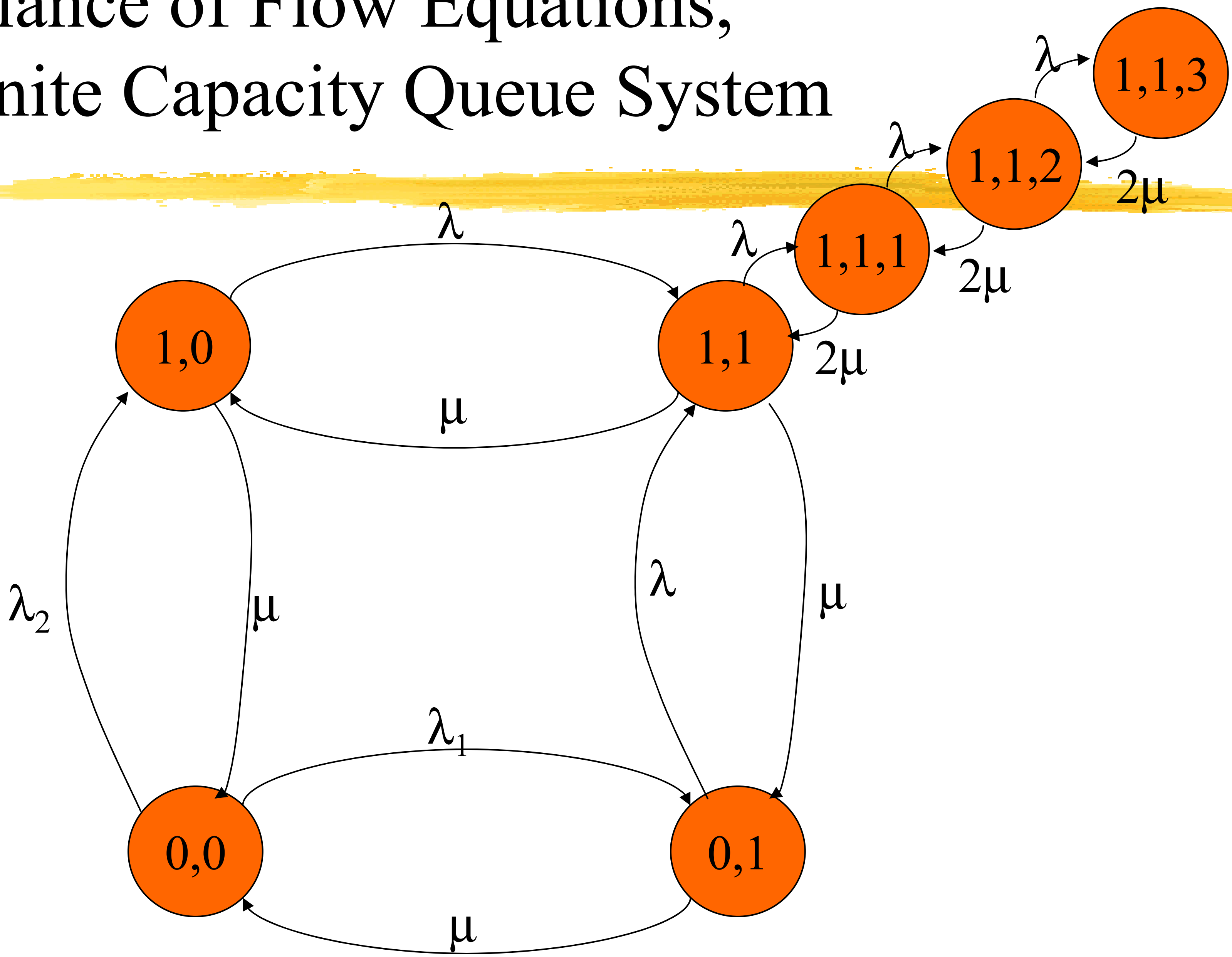
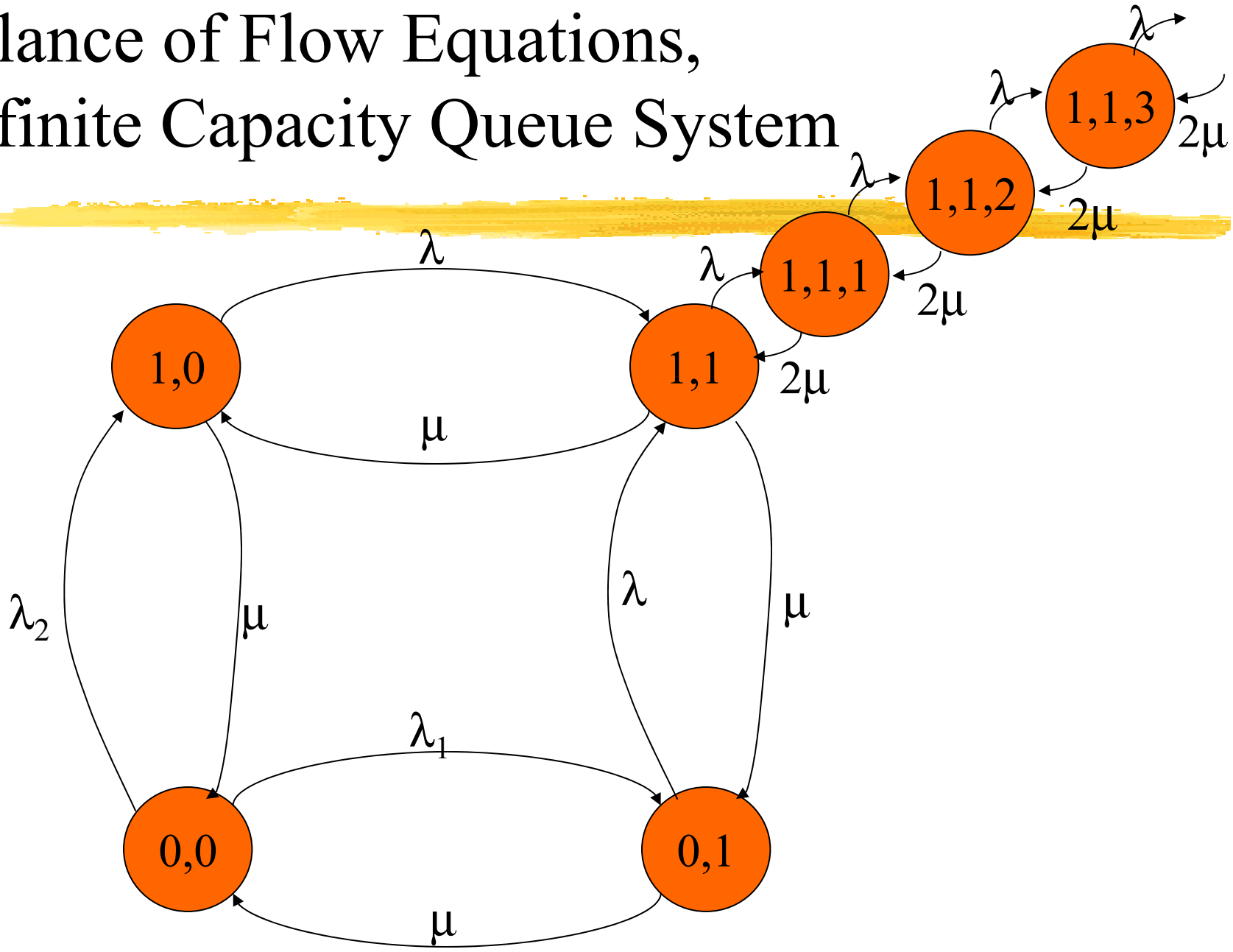
Balance of Flow Equations, Loss System



Balance of Flow Equations, Finite Capacity Queue System



Balance of Flow Equations, Infinite Capacity Queue System



Balance of Flow Equations, Loss System

$$P_{00}(\lambda_1 + \lambda_2) = P_{01}\mu + P_{10}\mu$$

$$P_{01}(\lambda + \mu) = P_{11}\mu + P_{00}\lambda_1$$

Etc.

$$P_{00} + P_{10} + P_{01} + P_{11} = 1$$

Workload and Imbalances

- $\rho_1 = W_1 = P_{01} + P_{11}$
- $\rho_2 = W_2 = P_{10} + P_{11}$
- *Workload Imbalance* = $\Delta W = |W_1 - W_2|$

To Obtain Travel Times, We Must Have Server Response Patterns

- f_{nj} = fraction of dispatches that are server *n* to response area *j*
- $T_n(\mathbf{C})$ = average time for server *n* to travel to a customer in region **C**
- $T(\mathbf{A})$ = average system-wide travel time, assuming that server 1's primary response area is region **A**.

Average System-Wide Travel Time

$$T(\mathbf{A}) = f_{11}T_1(\mathbf{A}) + f_{22}T_2(\mathbf{B\text{-}A}) + f_{12}T_1(\mathbf{B\text{-}A}) + f_{21}T_2(\mathbf{A})$$

Average System-Wide Travel Time

$$T(\mathbf{A}) = f_{11} T_1(\mathbf{A}) + f_{22} T_2(\mathbf{B\text{-}A}) + f_{12} T_1(\mathbf{B\text{-}A}) + f_{21} T_2(\mathbf{A})$$

Queueing

Average System-Wide Travel Time

$$T(\mathbf{A}) = f_{11} T_1(\mathbf{A}) + f_{22} T_2(\mathbf{B\text{-}A}) + f_{12} T_1(\mathbf{B\text{-}A}) + f_{21} T_2(\mathbf{A})$$

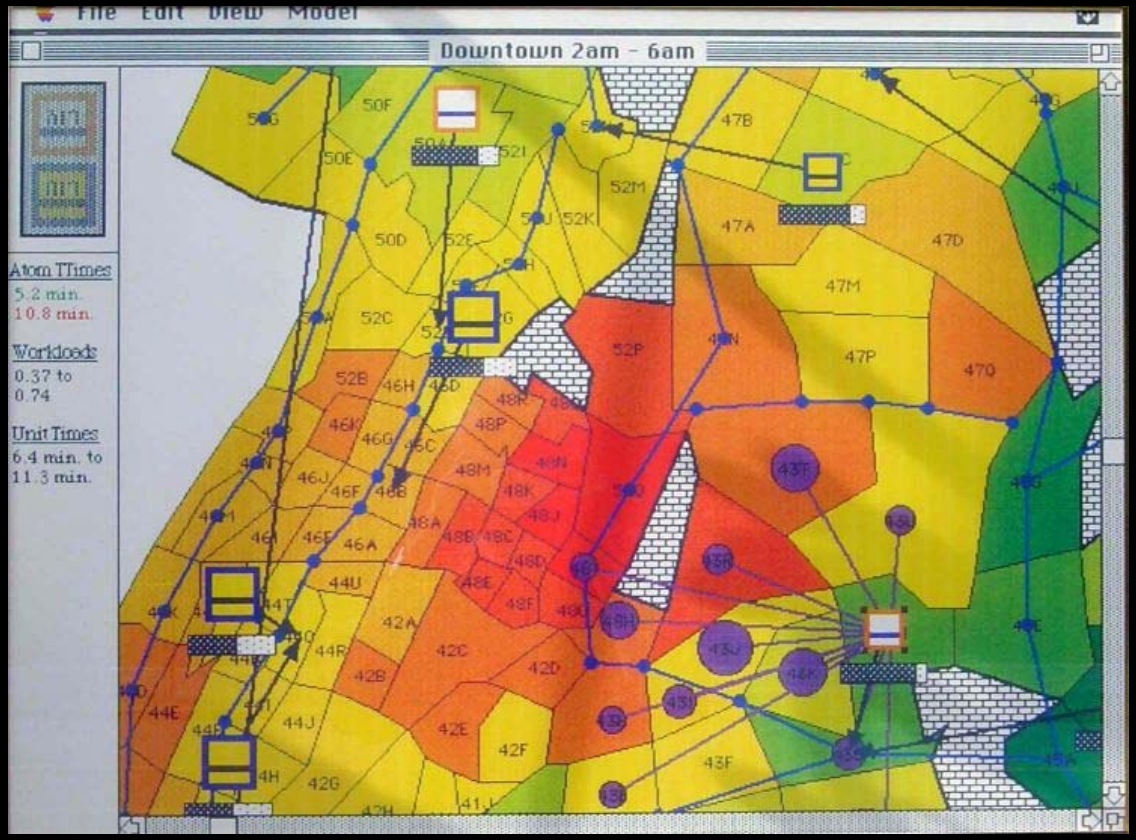
Geometry

How do we obtain the f_{nj}'s?

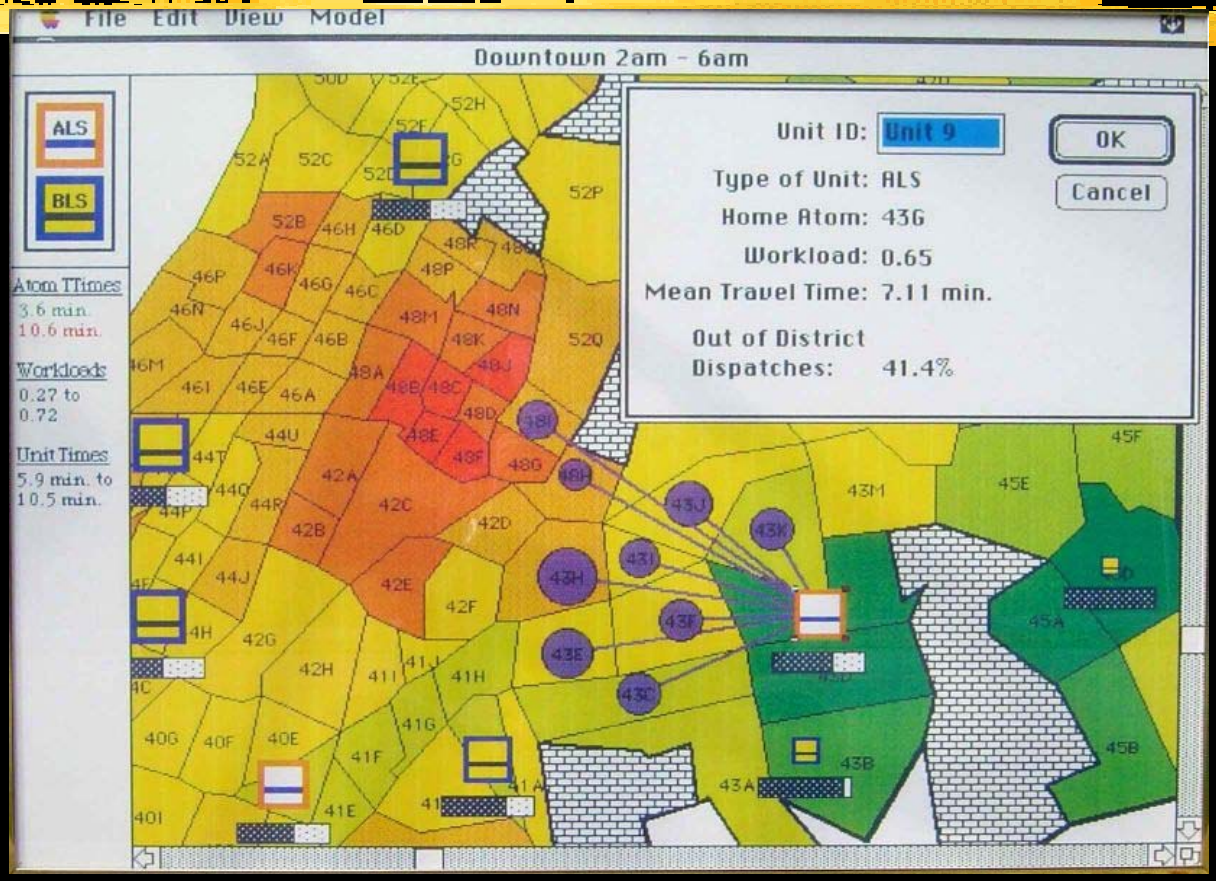
- Consider a long time interval T
- f_{12}=(# requests that assign unit 1 to area 2)/ (total # requests answered)
- Total # requests answered = $(1-P_{11})\lambda T$
- Average # requests that are "server 1 to area 2" is $\lambda_2 T P_{10}$. Why?
- Therefore $f_{12} = (\lambda_2 T P_{10} / [1-P_{11}]\lambda T) = \{\lambda_2/(1-P_{11})\lambda)\} P_{10}$

How do we generalize this to *N* servers?

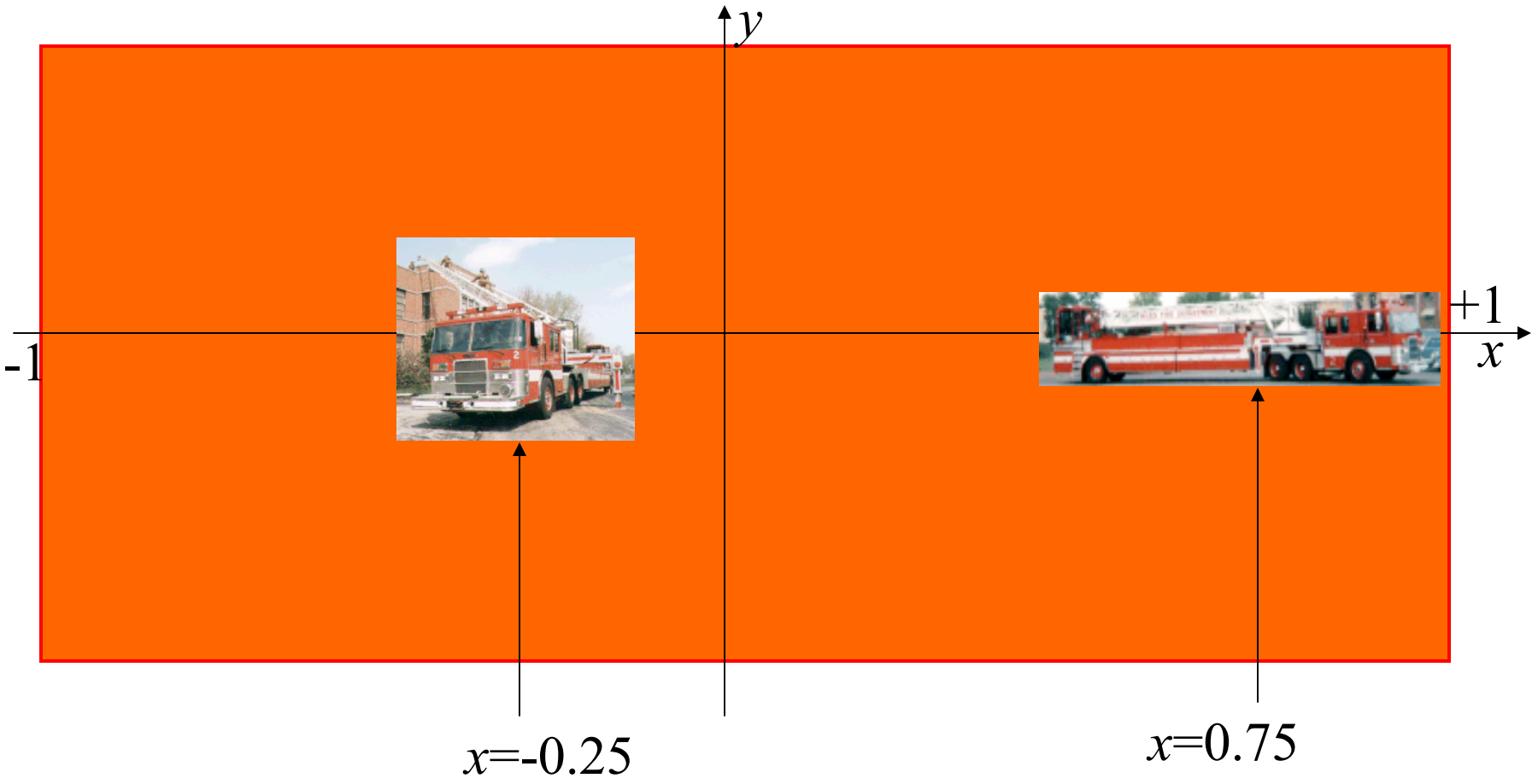
New York City EMS Hypercube



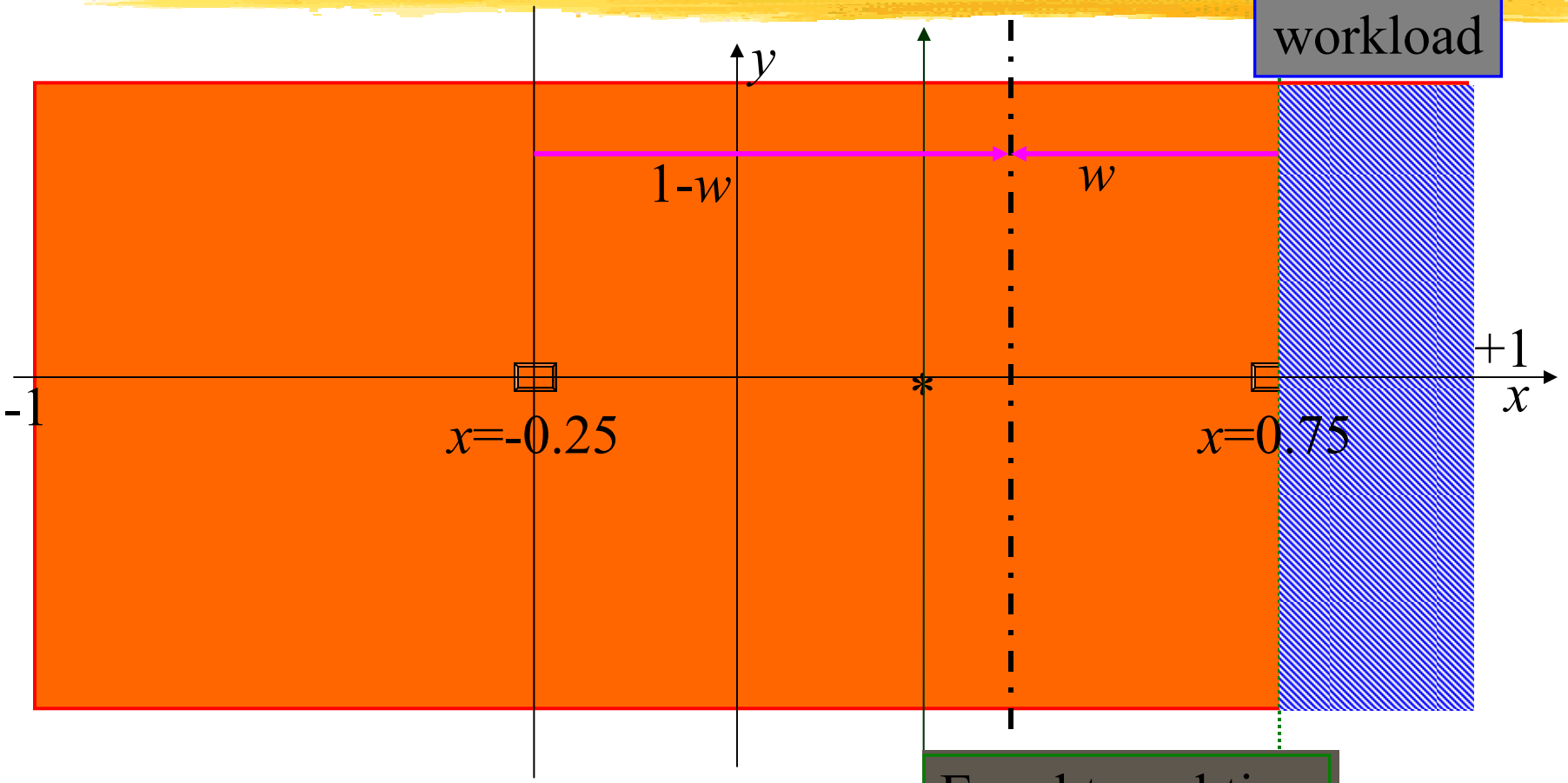
New York City EMS Hypercube



Rectangular City Example



Rectangular City Example



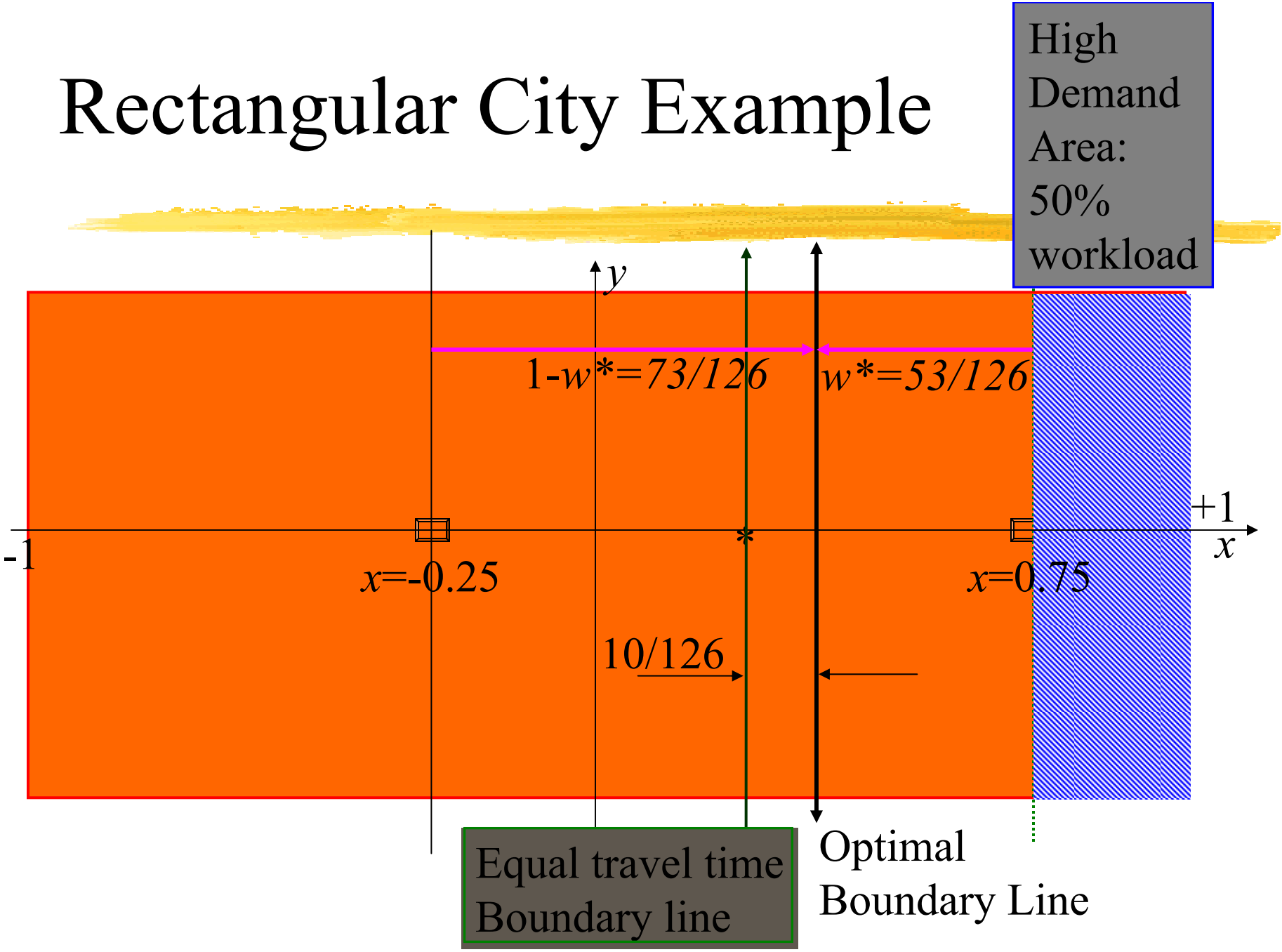
High Demand Area: 50% workload

Equal travel time Boundary line

Optimal Districting

- “Dispatch the closest available server’ is often not optimal, where ‘optimal’ implies minimizing mean travel time
- May not be good for reducing workload imbalance either
- With numerical example in book, the optimal boundary line is shifted to the right by 10/126 miles.

Rectangular City Example



Boundary Line Comparison

- Equal travel time boundary line
  - $T(A_{w=1/2})=0.46246$
  - $\Delta W = 0.05236$
- Optimal boundary line
  - $T(A_{w^*})=0.46166$
  - $\Delta W = 0.04405$

Two server Loss Model: Boundary Line Result

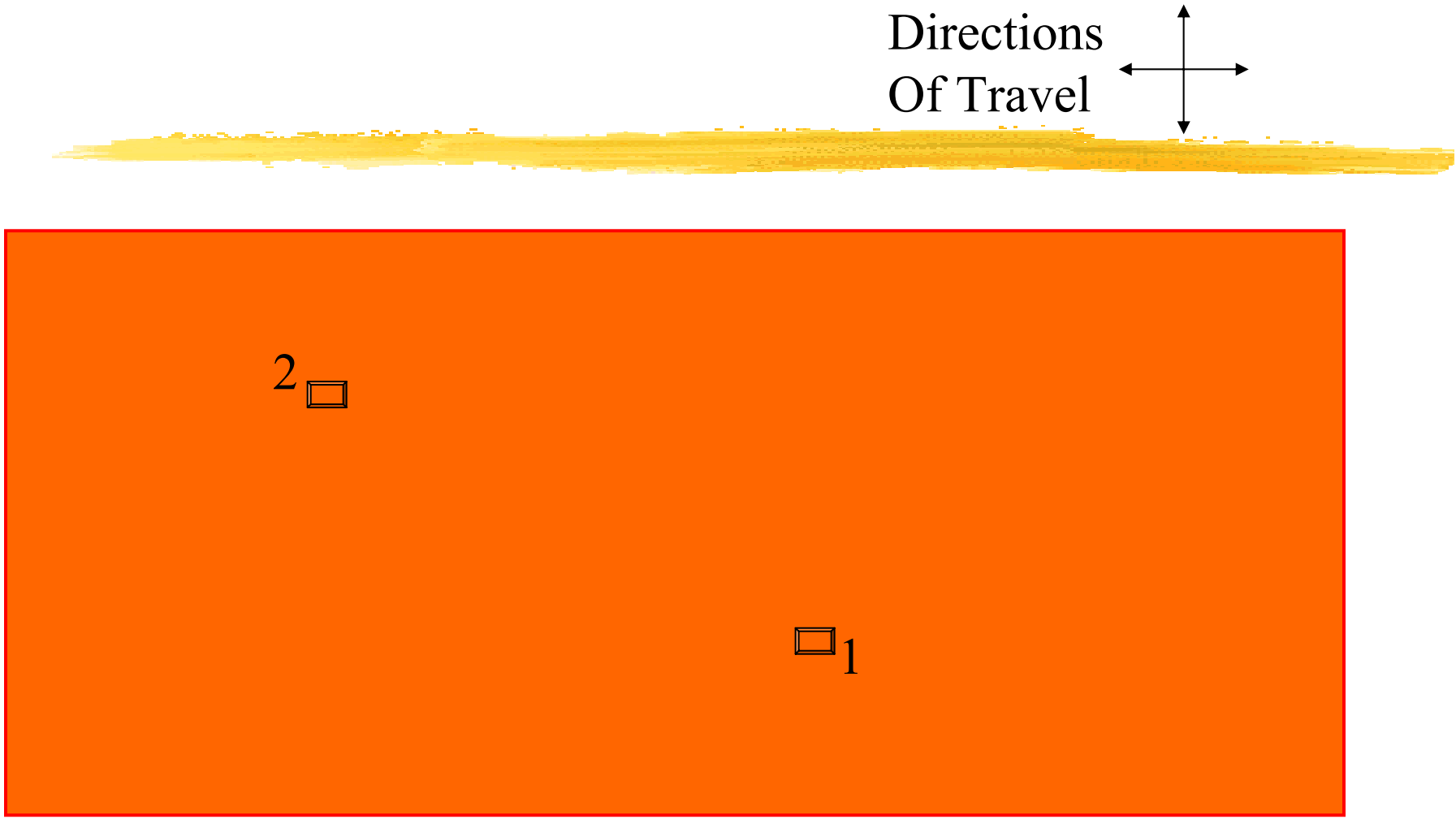
- **To minimize mean city-wide mean travel time:**
- ***The optimal partitioning consists of a set of points within the region that is a constant travel time s_0 closer to facility 1 than to facility 2.*** (Carter, Chaiken, Ignall, 1972)
- Does our rectangular city example work for this?

S_0: Optimal Partitioning

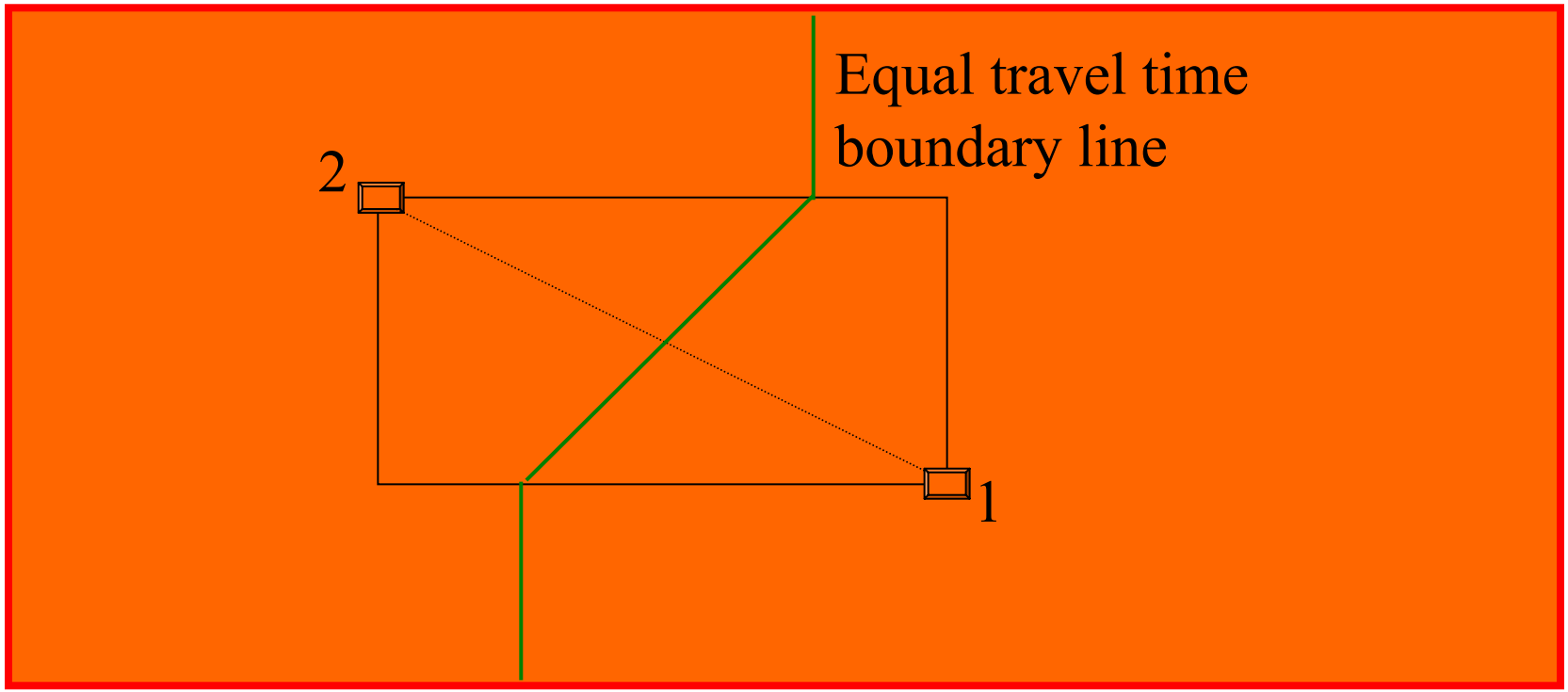
$$\alpha \equiv \lambda / 2\mu$$

$$\mu_1 = \mu_2$$

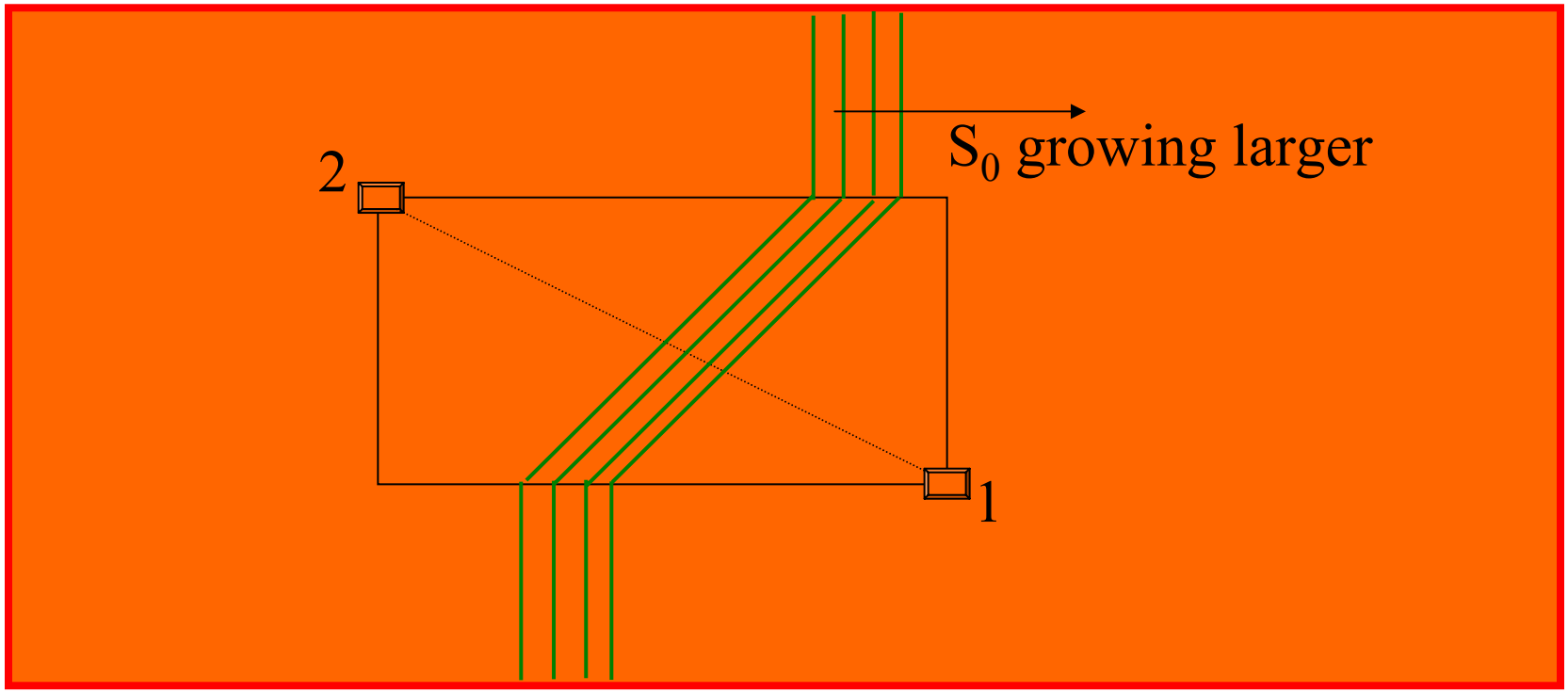
$$S_0 = [2\alpha /(2\alpha + 1)]\{T_2(B) - T_1(B)\}$$


Directions
Of Travel
2
1

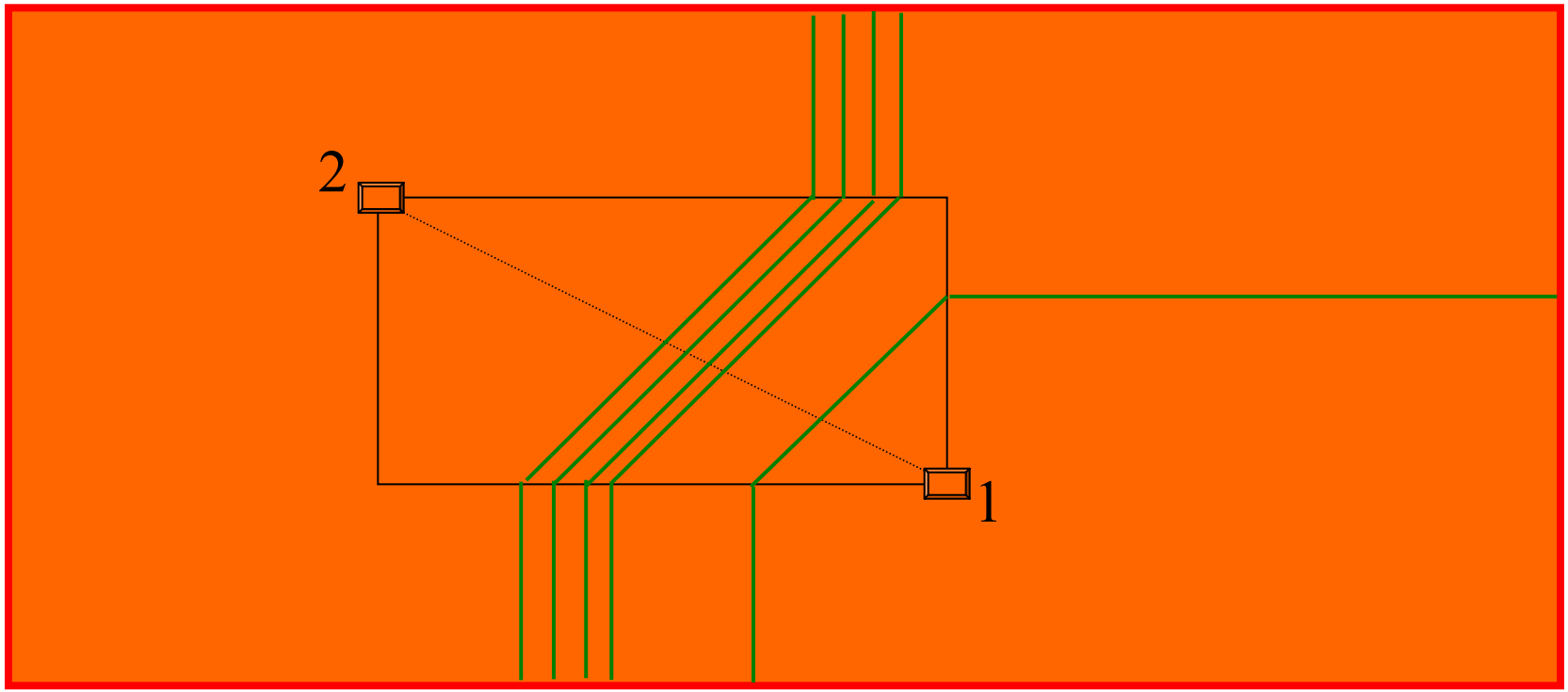
Directions
Of Travel



Directions
Of Travel

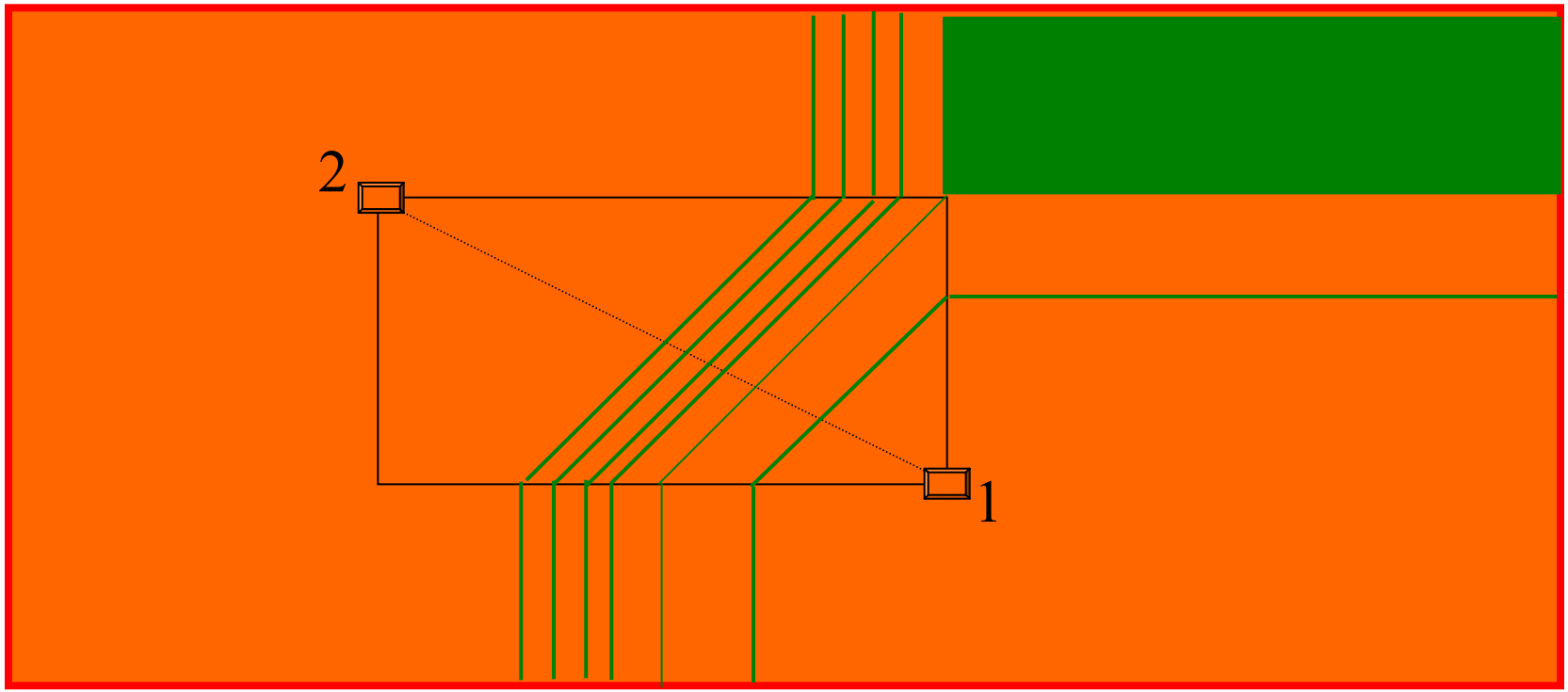


Directions
Of Travel



And what about the corner case?

Directions
Of Travel



The S_0 result is general

- Works for discrete grid
- One way streets
- General transportation network

- Rick Jarvis, in an MIT Ph.D. thesis, generalized this to *N* servers