

6.231 DYNAMIC PROGRAMMING

LECTURE 18

LECTURE OUTLINE

- Stochastic shortest path problems
- Policy iteration
- Linear programming
- Discounted problems

STOCHASTIC SHORTEST PATH PROBLEMS

- Assume finite-state system: States $1, \dots, n$ and special cost-free termination state t
 - Transition probabilities $p_{ij}(u)$
 - Control constraints $u \in U(i)$
 - Cost of policy $\pi = \{\mu_0, \mu_1, \dots\}$ is

$$J_\pi(i) = \lim_{N \rightarrow \infty} E \left\{ \sum_{k=0}^{N-1} g(x_k, \mu_k(x_k)) \mid x_0 = i \right\}$$

- Optimal policy if $J_\pi(i) = J^*(i)$ for all i .
- Special notation: For stationary policies $\pi = \{\mu, \mu, \dots\}$, we use $J_\mu(i)$ in place of $J_\pi(i)$.
- Assumption: There exists integer m such that for every policy and initial state, there is positive probability that the termination state will be reached after no more than m stages; for all π , we have

$$\rho_\pi = \max_{i=1, \dots, n} P\{x_m \neq t \mid x_0 = i, \pi\} < 1$$

MAIN RESULT

- Given any initial conditions $J_0(1), \dots, J_0(n)$, the sequence $J_k(i)$ generated by the DP iteration

$$J_{k+1}(i) = \min_{u \in U(i)} \left[g(i, u) + \sum_{j=1}^n p_{ij}(u) J_k(j) \right], \quad \forall i$$

converges to the optimal cost $J^*(i)$ for each i .

- Bellman's equation has $J^*(i)$ as unique solution:

$$J^*(i) = \min_{u \in U(i)} \left[g(i, u) + \sum_{j=1}^n p_{ij}(u) J^*(j) \right], \quad \forall i$$

- A stationary policy μ is optimal if and only if for every state i , $\mu(i)$ attains the minimum in Bellman's equation.

- Key proof idea: The “tail” of the cost series,

$$\sum_{k=mK}^{\infty} E \{ g(x_k, \mu_k(x_k)) \}$$

vanishes as K increases to ∞ .

BELLMAN'S EQUATION FOR A SINGLE POLICY

- Consider a stationary policy μ
- $J_\mu(i)$, $i = 1, \dots, n$, are the unique solution of the linear system of n equations

$$J_\mu(i) = g(i, \mu(i)) + \sum_{j=1}^n p_{ij}(\mu(i)) J_\mu(j), \quad \forall i = 1, \dots, n$$

- Proof: This is just Bellman's equation for a modified/restricted problem where there is only one policy, the stationary policy μ , i.e., the control constraint set at state i is $\tilde{U}(i) = \{\mu(i)\}$
- The equation provides a way to compute $J_\mu(i)$, $i = 1, \dots, n$, but the computation is substantial for large n [$O(n^3)$]

POLICY ITERATION

- It generates a sequence μ^1, μ^2, \dots of stationary policies, starting with any stationary policy μ^0 .
- At the typical iteration, given μ^k , we perform a *policy evaluation step*, that computes the $J_{\mu^k}(i)$ as the solution of the (linear) system of equations

$$J(i) = g(i, \mu^k(i)) + \sum_{j=1}^n p_{ij}(\mu^k(i)) J(j), \quad i = 1, \dots, n,$$

in the n unknowns $J(1), \dots, J(n)$. We then perform a *policy improvement step*, which computes a new policy μ^{k+1} as

$$\mu^{k+1}(i) = \arg \min_{u \in U(i)} \left[g(i, u) + \sum_{j=1}^n p_{ij}(u) J_{\mu^k}(j) \right], \quad \forall i$$

- The algorithm stops when $J_{\mu^k}(i) = J_{\mu^{k+1}}(i)$ for all i
- Note the connection with the rollout algorithm, which is just a single policy iteration

JUSTIFICATION OF POLICY ITERATION

- We can show that $J_{\mu^{k+1}}(i) \leq J_{\mu^k}(i)$ for all i, k
- Fix k and consider the sequence generated by

$$J_{N+1}(i) = g(i, \mu^{k+1}(i)) + \sum_{j=1}^n p_{ij}(\mu^{k+1}(i)) J_N(j)$$

where $J_0(i) = J_{\mu^k}(i)$. We have

$$\begin{aligned} J_0(i) &= g(i, \mu^k(i)) + \sum_{j=1}^n p_{ij}(\mu^k(i)) J_0(j) \\ &\geq g(i, \mu^{k+1}(i)) + \sum_{j=1}^n p_{ij}(\mu^{k+1}(i)) J_0(j) = J_1(i) \end{aligned}$$

Using the monotonicity property of DP,

$$J_0(i) \geq J_1(i) \geq \dots \geq J_N(i) \geq J_{N+1}(i) \geq \dots, \quad \forall i$$

Since $J_N(i) \rightarrow J_{\mu^{k+1}}(i)$ as $N \rightarrow \infty$, we obtain $J_{\mu^k}(i) = J_0(i) \geq J_{\mu^{k+1}}(i)$ for all i . Also if $J_{\mu^k}(i) = J_{\mu^{k+1}}(i)$ for all i , J_{μ^k} solves Bellman's equation and is therefore equal to J^*

- A policy cannot be repeated, there are finitely many stationary policies, so the algorithm terminates with an optimal policy

LINEAR PROGRAMMING

- We claim that J^* is the “largest” J that satisfies the constraint

$$J(i) \leq g(i, u) + \sum_{j=1}^n p_{ij}(u) J(j), \quad (1)$$

for all $i = 1, \dots, n$ and $u \in U(i)$.

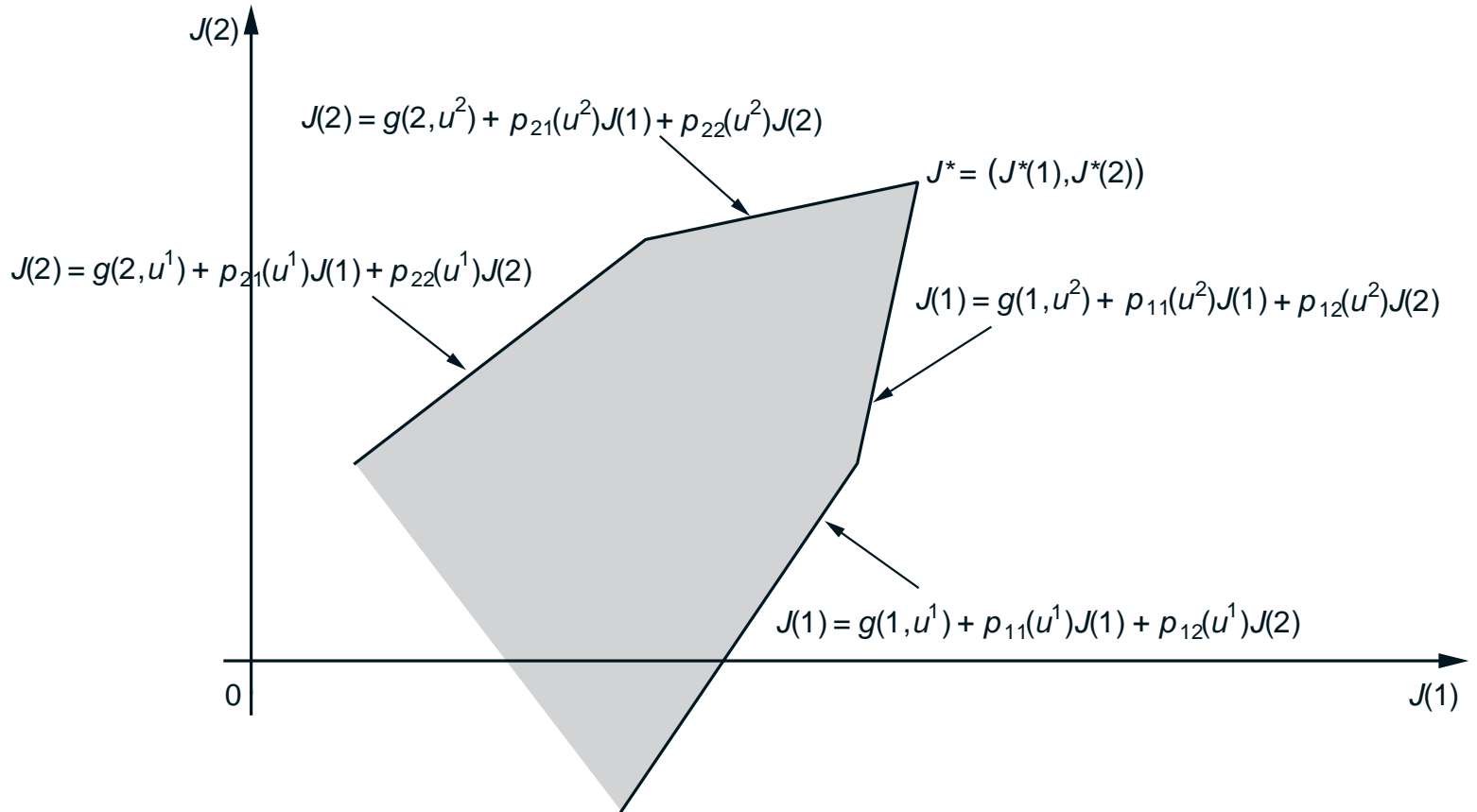
- Proof: If we use value iteration to generate a sequence of vectors $J_k = (J_k(1), \dots, J_k(n))$ starting with a J_0 such that

$$J_0(i) \leq \min_{u \in U(i)} \left[g(i, u) + \sum_{j=1}^n p_{ij}(u) J_0(j) \right], \quad \forall i$$

Then, $J_k(i) \leq J_{k+1}(i)$ for all k and i (monotonicity of DP) and $J_k \rightarrow J^*$, so that $J_0(i) \leq J^*(i)$ for all i .

- So $J^* = (J^*(1), \dots, J^*(n))$ is the solution of the linear program of maximizing $\sum_{i=1}^n J(i)$ subject to the constraint (1).

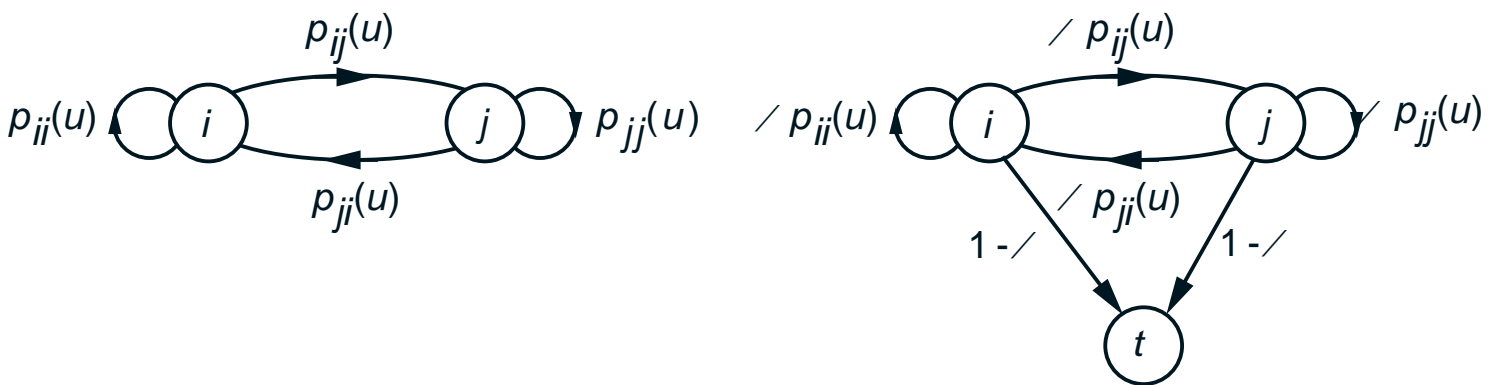
LINEAR PROGRAMMING (CONTINUED)



- Drawback: For large n the dimension of this program is very large. Furthermore, the number of constraints is equal to the number of state-control pairs.

DISCOUNTED PROBLEMS

- Assume a discount factor $\alpha < 1$.
- Conversion to an SSP problem.



- Value iteration converges to J^* for all initial J_0 :

$$J_{k+1}(i) = \min_{u \in U(i)} \left[g(i, u) + \alpha \sum_{j=1}^n p_{ij}(u) J_k(j) \right], \forall i$$

- J^* is the unique solution of Bellman's equation:

$$J^*(i) = \min_{u \in U(i)} \left[g(i, u) + \alpha \sum_{j=1}^n p_{ij}(u) J^*(j) \right], \forall i$$

DISCOUNTED PROBLEMS (CONTINUED)

- Policy iteration converges finitely to an optimal, and linear programming works.
- Example: Asset selling over an infinite horizon. If accepted, the offer x_k of period k , is invested at a rate of interest r .
- By depreciating the sale amount to period 0 dollars, we view $(1 + r)^{-k} x_k$ as the reward for selling the asset in period k at a price x_k , where $r > 0$ is the rate of interest. So the discount factor is $\alpha = 1/(1 + r)$.
- J^* is the unique solution of Bellman's equation

$$J^*(x) = \max \left[x, \frac{E\{J^*(w)\}}{1 + r} \right].$$

- An optimal policy is to sell if and only if the current offer x_k is greater than or equal to $\bar{\alpha}$, where

$$\bar{\alpha} = \frac{E\{J^*(w)\}}{1 + r}.$$