

# Game Theory is Evolving

MIT 14.126, Fall 2001

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# Our Topics in the Course

## ◆ Classical Topics

- Choice under uncertainty
- Cooperative games
  - ◆ Values
  - ◆ 2-player bargaining
  - ◆ Core and related
  - ◆ Cores of market games
- Non-cooperative games
  - ◆ Fixed points/equilibrium
  - ◆ Refinements
  - ◆ (Repeated games)

## ◆ Newer Topics

- "Nash program" (noncoop foundations for coop games)
- Cheap talk
- Experiments
- Foundations of Noncoop Games
  - ◆ Epistemic
  - ◆ Learning
  - ◆ Evolutionary
- Mechanism design
- Other economic applications
  - ◆ Commitment, signaling
  - ◆ Comparative statics

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# Two Game Formulations

## The cooperative concept

- Perfect communication
- Perfect contract enforcement

## ◆ Formulation $(N, v)$

- $N$  is a set of players
- TU case:  $v: 2^N \rightarrow \mathbb{R}$
- NTU case:  $v(S): 2^N \rightarrow \mathbb{R}^S$

- ◆ Solution(s): (set of) value allocations  $\pi \in \mathbb{R}^N$

## The non-cooperative concept

- No communication
- No contract enforcement

## ◆ Formulation $(N, S, \pi)$

- $N$  is a set of players
- $x \in S = S_1 \times \dots \times S_N$  is a strategy profile
- $\pi_n(x) = n$ 's payoff

- ◆ Solution(s): (set of) strategy profiles  $x \in S$

# Example: Linear Duopoly

## ◆ Non-cooperative

- Players: 1 and 2
- Strategies: Output levels  $x_1$  and  $x_2$
- Payoffs:  $x_i[1 - (x_1 + x_2) - c_i]$
- Nash Equilibrium:  
 $x_1 = (1 + c_2 - 2c_1)/3$   
 $x_2 = (1 + c_1 - 2c_2)/3$

## ◆ Cooperative TU

$$v(1) = v(2) = 0?$$
$$v(1,2) = \max_x x(1 - x - \min(c_1, c_2))$$

## ◆ Cooperative NTU

$$v(1) = v(2) = 0?$$
$$v(1,2) = \{(\pi_1, \pi_2) : (\exists x_1, x_2 \in [0,1])$$
$$\pi_1 = x_1(1 - x_1 - x_2 - c_1) \ \&$$
$$\pi_2 = x_2(1 - x_1 - x_2 - c_2)\}$$

## ◆ Solution?

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## Cooperative Solutions

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## Solutions: Points & Sets

- ◆ Cooperative “one point” solutions usually attempt to extend a symmetric solution to asymmetric games.
  - “Nash bargaining solution”
  - “Shapley value”
  - In the TU Cournot game, the solution might be  $\pi_j = v(i) + 1/2[v(1,2) - v(i)]$  with  $v(i) = 0$  or Cournot profit.
- ◆ Others generalize the indeterminacy of the bargaining outcome and identify only what is “blocked.”
  - In the NTU Cournot game, the “core” consists of all the efficient, individually rational outcomes.

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## The Shapley Value $\varphi_n(N, v)$

### ◆ Four axioms for TU games

- Efficiency:  $\sum \varphi_n(N, v) = v(N)$
- Null player: If  $v(S \cup \{n\}) = v(S)$  then  $\varphi_n(N, v) = 0$ .
- Symmetry: Permutations don't matter
- Additivity:  $\varphi(N, v+w) = \varphi(N, v) + \varphi(N, w)$

### ◆ Analysis:

- The games  $(N, v)$  (with  $N$  fixed) form a linear space.
- Let  $\chi_S(T) = 1$  if  $S \subset T$  and  $\chi_S(T) = 0$  otherwise.
- Axiom 1-3:  $\varphi_n(\alpha \chi_S) = \alpha / |S|$  if  $n \in S$  and  $\varphi_n(\alpha \chi_S) = 0$  if  $n \notin S$ .
- Add Axiom 4:  $\varphi$  is a linear operator

## Shapley's Theorem

### ◆ Notation:

- $\Pi$  = set of permutations of  $N$ , typical element  $\pi$ , mapping elements of  $N$  onto  $\{1, \dots, |N|\}$ .
- $S_{i\pi} = \{j: \pi_j \leq \pi_i\}$

### ◆ Theorem.

$$\varphi_n(N, v) = \frac{1}{|N|!} \sum_{\pi \in \Pi} v(S_{i\pi}) - v(S_{i\pi} \setminus i)$$

- ◆ Proof. The given  $\varphi$  satisfies the axioms by inspection. Since the  $\chi_S$  games form a basis, there is a unique linear operator that does so. QED

# Shapley Interpretations

- ◆ Power, for example in voting games.
- ◆ Fairness, for example in cost allocations.
- ◆ Extension: Aumann-Shapley pricing
  - Cornell long-distance phone cost allocation
  - Game has  $K$  types of players with mass  $a_k$  of type  $k$ .
  - "Diagonal formula":

$$\varphi_k(N, v) = \int v_k(sa_1, \dots, sa_K) ds$$

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# Crossover Ideas

Connecting Cooperative and Non-cooperative game ideas

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# Crossover Ideas, 1

- ◆ Cheap Talk
  - What happens to non-cooperative games when we add a stage of message exchange?
  - In the following game, suppose
    - ◆ Payoff unit is \$10,000s
    - ◆ Row player can send a message...

7,6	8,5
0,0	9,9

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# Crossover Ideas, 2

- ◆ The "Nash Program"
  - Non-cooperative foundations for cooperative solutions
  - 2-player bargaining problem
    - ◆ Nash's bargaining solution
    - ◆ Nash demand game
    - ◆ Stahl-Rubinstein "alternating offers" model
  - Exchange economies
    - ◆ Competitive equilibrium (core)
    - ◆ Shapley-Shubik game
    - ◆ Auction and bidding games

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# Nash Demand Bargaining

- ◆ Let  $v(1)=v(2)=0$ ,  $v(12)=1$ .
  - The "core" is the set of payoff vectors  $\pi$  that are:
    - ◆ "Efficient":  $\sum_{n \in N} \pi_n = v(N)$
    - ◆ "Unblocked":  $(\forall S \subset N) \sum_{n \in S} \pi_n \geq v(S)$
  - In this case,  $\text{Core}(N, v) = \{(\pi_1, \pi_2) : \pi_1, \pi_2 \geq 0, \pi_1 + \pi_2 = 1\}$ .
- ◆ Nash demand game.  $S_1 = S_2 = [0, 1]$ .
  - $\pi_n = \begin{cases} x_n & \text{if } x_1 + x_2 \leq 1 \\ 0 & \text{otherwise} \end{cases}$
  - "Perfect" equilibria coincide with the core.

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# Alternating Offer Bargaining

- ◆ Players 1 and 2 take turn making offers.
- ◆ If agreement  $(x_1, x_2)$  is reached at time  $t$ , payoffs are  $\delta^t(x_1, x_2)$ .
- ◆ Unique subgame perfect equilibrium has payoffs of "almost" the  $(.5, .5)$  Nash solution:

$$\pi_n = \begin{cases} \frac{1}{1+\delta} & \text{if } n \text{ moves first} \\ \frac{\delta}{1+\delta} & \text{otherwise} \end{cases}$$

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# Noncooperative Theory

Rethinking equilibrium

# What are we trying to model?

- ◆ Animal behavior:
  - evolutionary stable strategies
- ◆ Learned behavior:
  - Reinforcement learning models
  - Self-confirming equilibrium
- ◆ Self-enforcing agreements:
  - Nash equilibrium & refinements
- ◆ Reflection among rational players
  - "Interactive epistemology"

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# Animal Behavior

- ◆ Hawk-Dove game: what behavior will evolve?

-1,-1	2,0
0,2	1,1

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# Human Experiments

## ◆ The "Ultimatum Game"

- Players: {1,2}
- Strategies
  - ◆ Player 1 makes an offer  $x, 1-x$ ;  $x \in [0,1]$
  - ◆ Player 2 sees it and says "yes" or "no"
- Payoffs
  - ◆ If 2 says "yes," payoffs are  $(x, 1-x)$ .
  - ◆ If 2 says "no," payoffs are  $(0,0)$ .

## ◆ Analysis

- Unique subgame perfect equilibrium outcome:  $(x=1, \text{yes})$ .
- It never happens that way!

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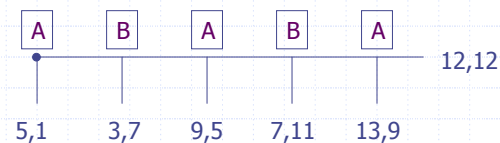
# Learning Models

- ◆ Behavior is established by some kind of adaptive learning
  - Kind #1: Players repeat strategies that were "successful" in "similar" past situations
  - Kind #2: Players forecast based on others' past play and, eventually, optimize accordingly.
    - ◆ Stationary points: "fulfilled expectations" equilibrium

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# Epistemic Analyses

- ◆ "What I do depends on what I expect."
- ◆ The "centipede game"
  - Is it rational below for A to grab \$5?
  - What should B believe if its probability one forecast of A's behavior is contradicted?



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# Forward Induction

- ◆ Are both subgame perfect equilibria reasonable as rational play?

