

Forward Induction, Signaling and Reputation

14.126 Game Theory

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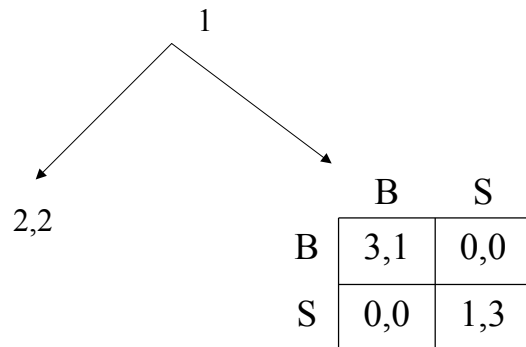
Muhamet Yildiz

Road Map

1. Forward Induction
2. Signaling games
 1. Sequential Equilibria
 2. Intuitive Criteria
3. Reputation
 1. Chain-store paradox, finitely repeated games
 2. Centipede game with incomplete information
 3. Finitely repeated entry-deterrence game with incomplete information.

Forward Induction

The Battle of the Sexes with outside options



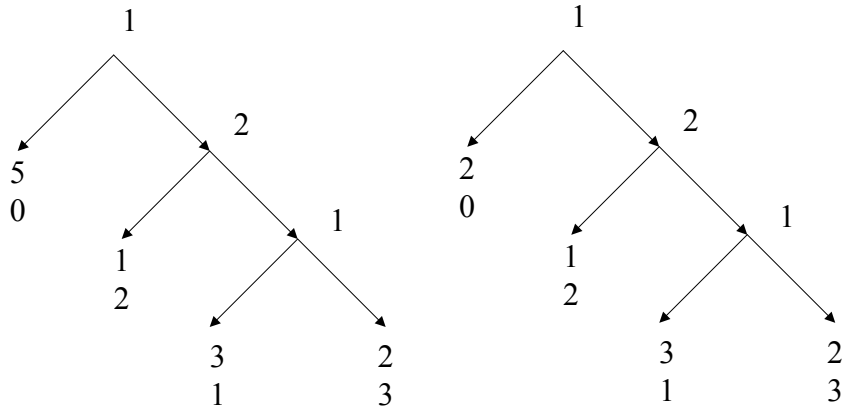
Forward Induction

- One ought to interpret the actions as outcomes of conscious choice even off the path.
- Intuitive criterion
- Mistaken theories

Strong belief in rationality

At any history of the game, each agent is assumed to be rational if possible. (That is, if there are two strategies s and s' of a player i that are consistent with a history of play, and if s is strictly dominated but s' is not, at this history no player j believes that i plays s .)

Examples



Burning Money

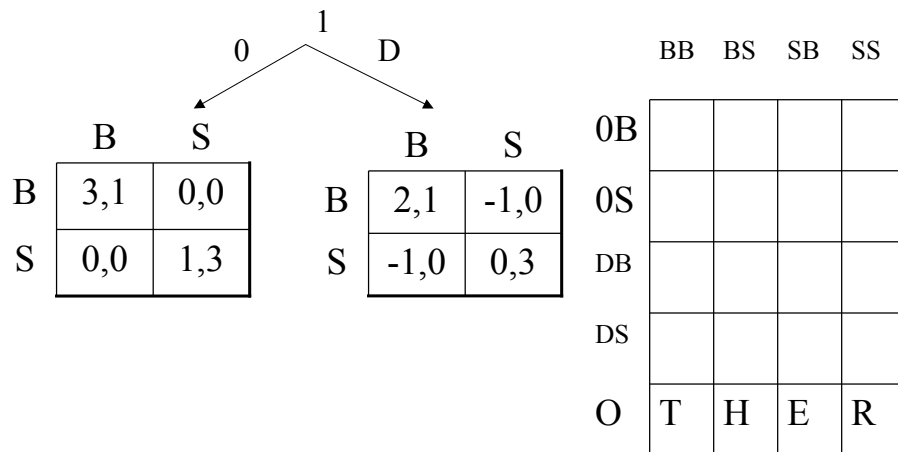


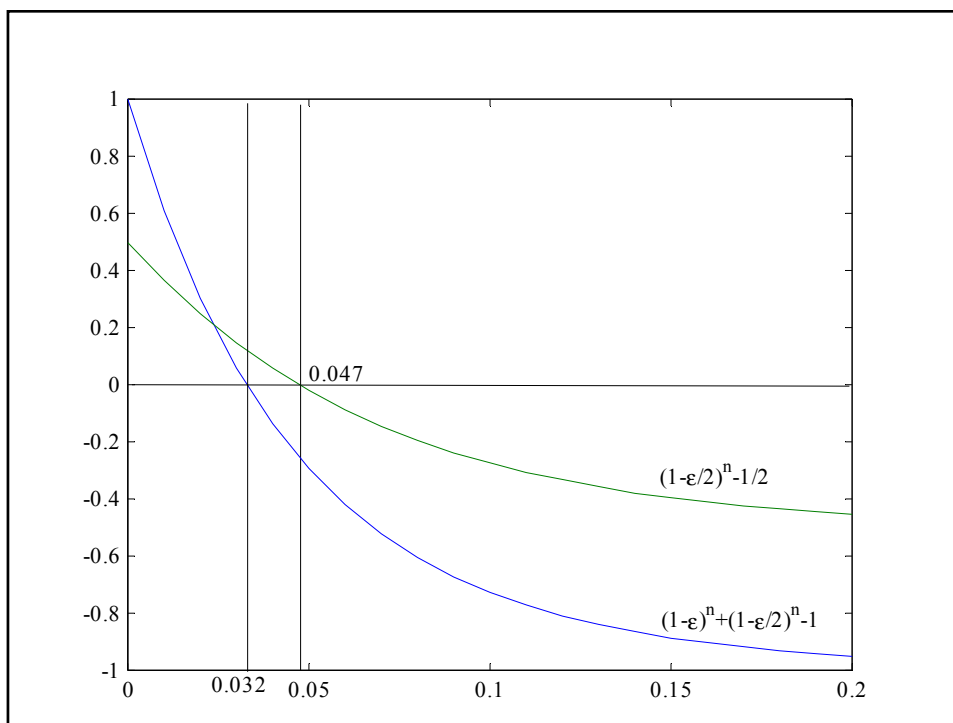
Table for the bidding game

$$U_i = 20(2 + 2\min_j \text{bid}_j - \text{bid}_i)$$

min bid \	1	2	3
1	60	-	-
2	40	80	-
3	20	60	100

Nash equilibria of bidding game

- 3 equilibria: s^1 = everybody plays 1; s^2 = everybody plays 2; s^3 = everybody plays 3.
- Assume each player trembles with probability $\epsilon < 1/2$, and plays each unintended strategy w.p. $\epsilon/2$, e.g., w.p. $\epsilon/2$, he thinks that such other equilibrium is to be played.
 - s^3 is an equilibrium iff
 - s^2 is an equilibrium iff
 - s^1 is an equilibrium iff



Bidding game with entry fee

Each player is first to decide whether to play the bidding game (E or X); if he plays, he is to pay a fee $p > 60$.

min \ Bid	1	2	3
1	60	-	-
2	40	80	-
3	20	60	100

For each $m = 1, 2, 3$, \exists SPE: (m, m, m) is played in the bidding game, and players play the game iff $20(2+m) \geq p$.

Forward induction: when $20(2+m) < p$, (E_m) is strictly dominated by (X_k) . After E, no player will assign positive probability to $\min \text{ bid} \leq m$. FI-Equilibria: (E_m, E_m, E_m) where $20(2+m) \geq p$.

What if an auction before the bidding game?

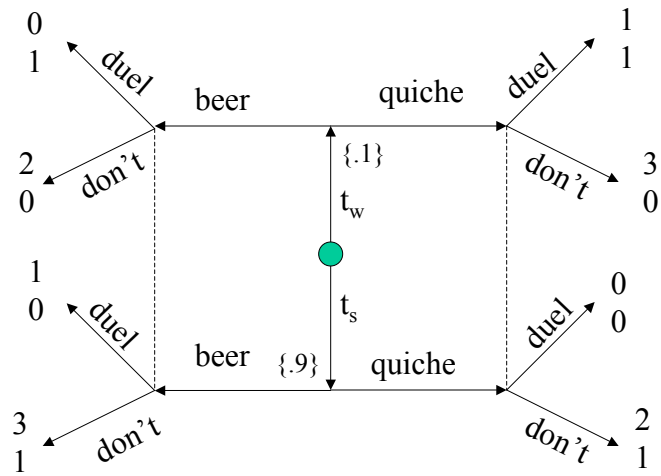
Signaling

Model

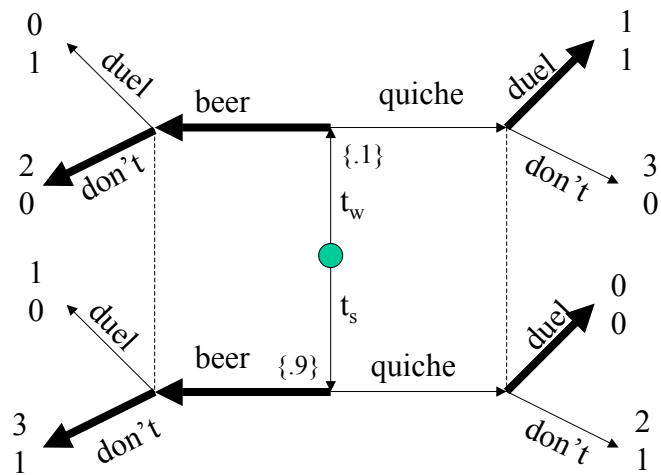
- Players: (S)ender, (R)eceiver
- 1. Nature selects t from T – the probability distribution is π ;
- 2. S observes t , and sends message m from a set M ;
- 3. R observes m – but not t – and takes an action a ;
- 4. S gets $U^S(t,m,a)$ and R gets $U^R(t,m,a)$.

This is common knowledge.

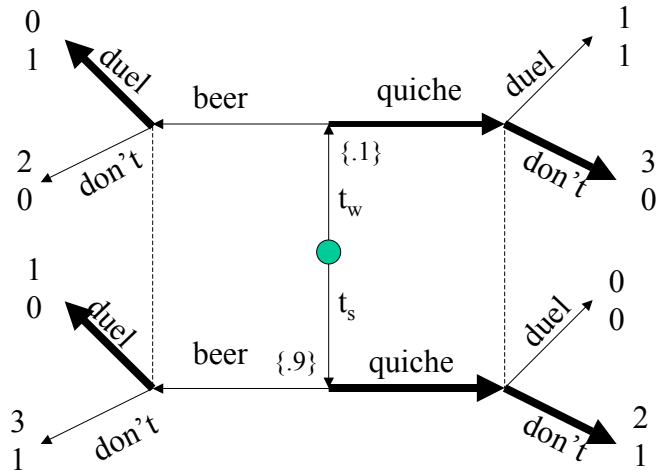
Beer – Quiche



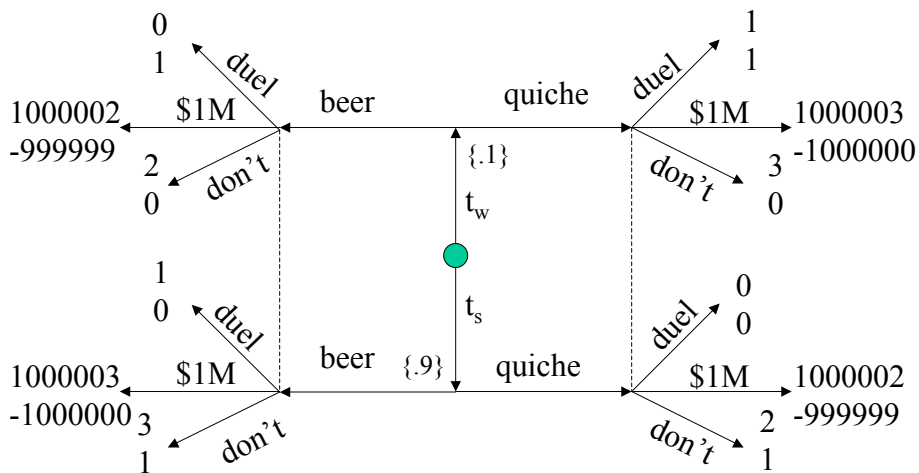
Good equilibrium



Bad equilibrium



Beer – Quiche – M



Cho -- Kreps

- $T(m); M(t); A(m)$
- Action set is finite;
- $\rho(m;t) =$ the probability that t sends m ;
- $\phi(a;m) =$ the probability that R chooses a , receiving m ;
- $BR(\mu, m) = \arg \max_{a \in A(m)} \sum_{t \in T(m)} U^R(t, m, a) \mu(t)$
- For subset I of T ,

$$BR(I, m) = \bigcup_{\{\mu: \mu(I)=1\}} BR(\mu, m)$$
- MBR

Sequential equilibrium

- Beliefs:

$$\mu(t | m) = \begin{cases} \frac{\pi(t)\rho(m;t)}{\sum_{t' \in T(m)} \pi(t')\rho(m;t')} & \text{if } \sum_{t' \in T(m)} \pi(t')\rho(m;t') > 0 \\ \text{something} & \text{otherwise} \end{cases}$$
- $\rho(m,t) > 0 \Rightarrow \sum_a U^S(t,m,a)\phi(a;m)$ is maximized at m ;
- $\phi(\cdot; m)$ is in $MBR(\mu(\cdot|m), m)$

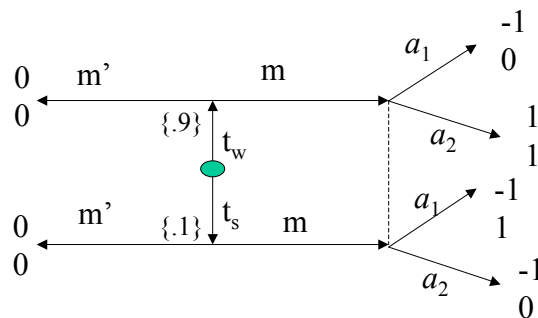
Testing an equilibrium

- $U^*(t)$ = expected utility of type t in equilibrium;
1. Pick a criterion, saying that particular out-of-equilibrium message (OEM) cannot be sent by some type t . Also, say that a will not be taken in response to m if a is not in $BR(T(m), m)$. Iterate. $[T^s(m)]$
 2. For each OEM m , consider all sequential equilibrium responses of R to m in the original game. Are all of these sequentially rational, given $T^s(m)$. If not, FAIL.

Dominance

- For any OEM m , eliminate t if $\exists m'$ s.t.

$$\min_{a \in A(m')} U^S(t, m', a) > \max_{a \in A(m)} U^S(t, m, a)$$



Equilibrium Dominance & Intuitive Criterion

Equilibrium Domination: \forall OEM m , eliminate (t,m) if

$$U^*(t) > \max_{a \in A(m)} U^S(t, m, a).$$

Intuitive Criterion: \forall OEM m , define

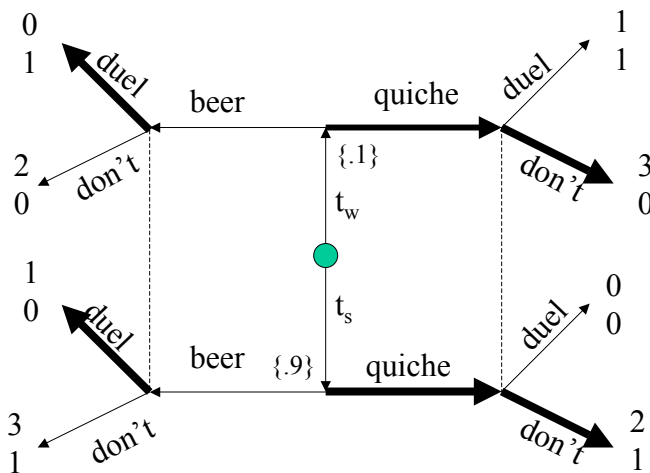
$$\tilde{T}(m) = \{t \mid U^*(t) > \max_{a \in BR(T(m),m)} U^S(t, m, a)\}$$

If $\exists (t',m)$ s.t.

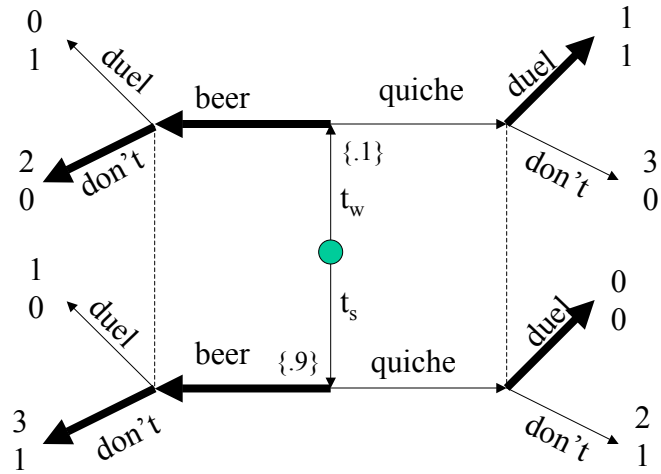
$$U^*(t') < \min_{a \in BR(T(m) \setminus \tilde{T}(m),m)} U^S(t', m, a),$$

then the equilibrium fails the Intuitive Criterion.

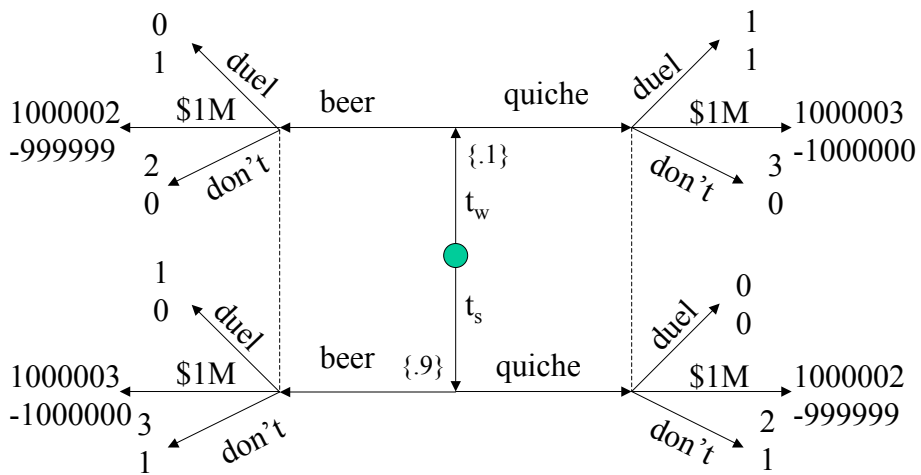
Bad equilibrium



Good equilibrium

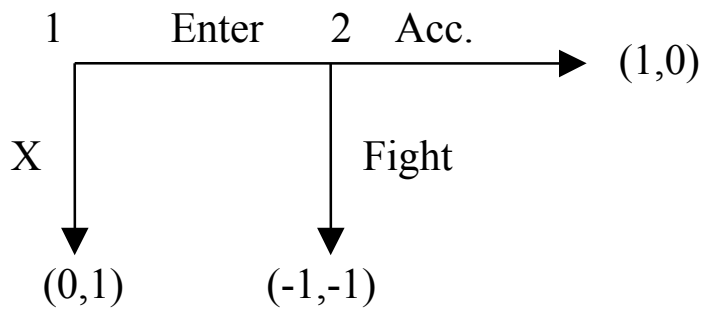


Beer – Quiche – M

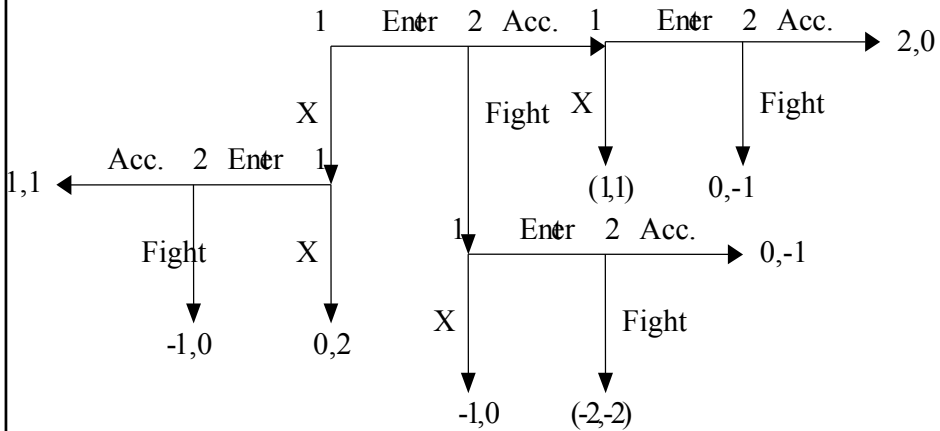


Reputation

Entry deterrence

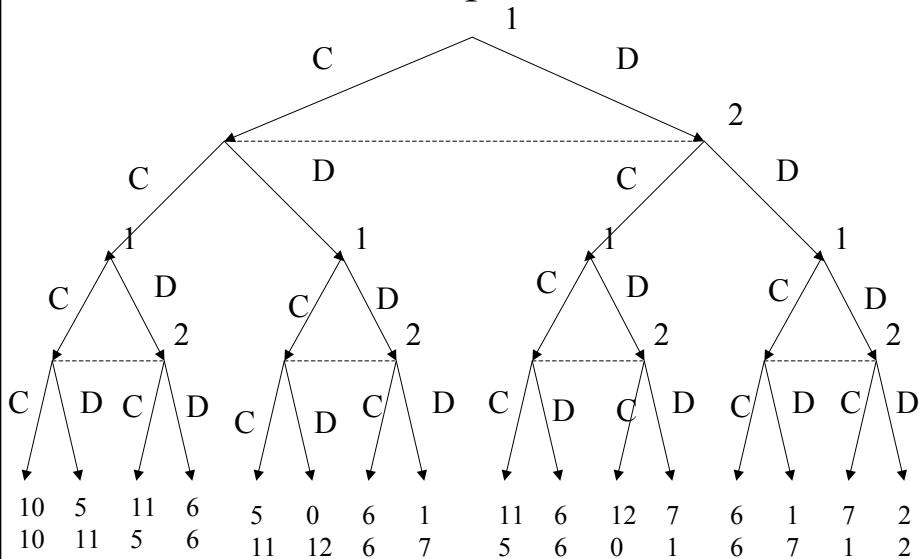


Entry deterrence, repeated twice, many times



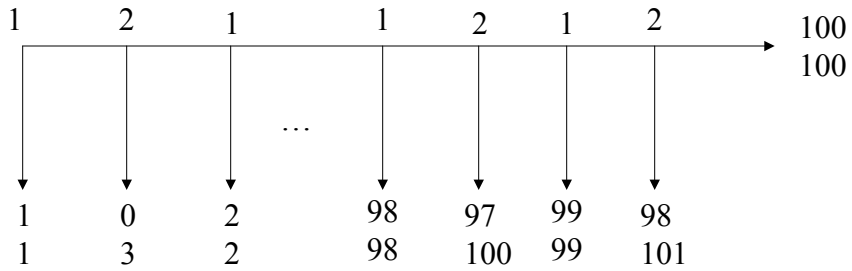
What would happen if repeated n times?

Twice-repeated PD

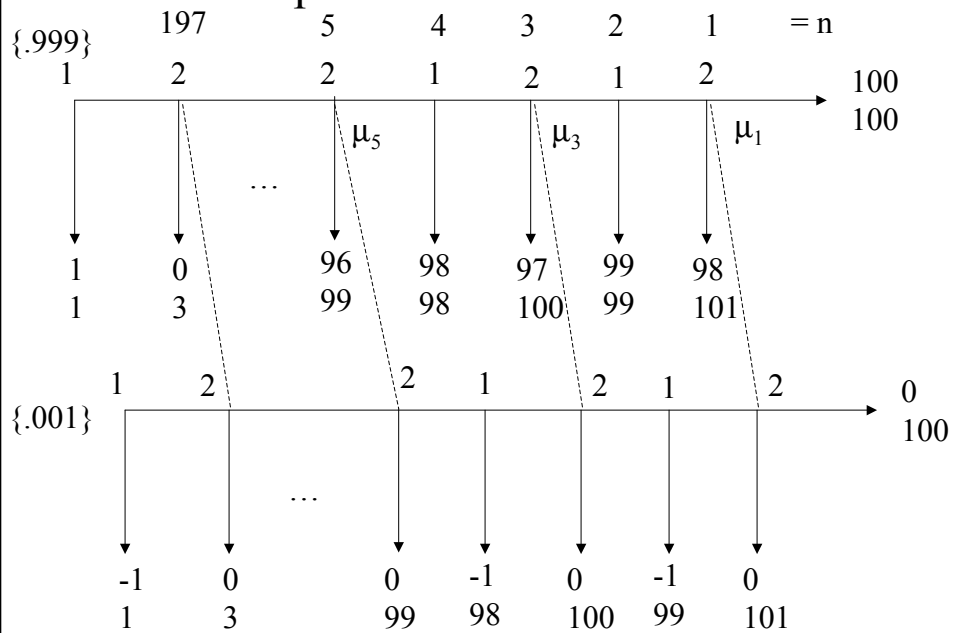


What would happen if $T = \{0,1,2,\dots,n\}$?

Centipede Game

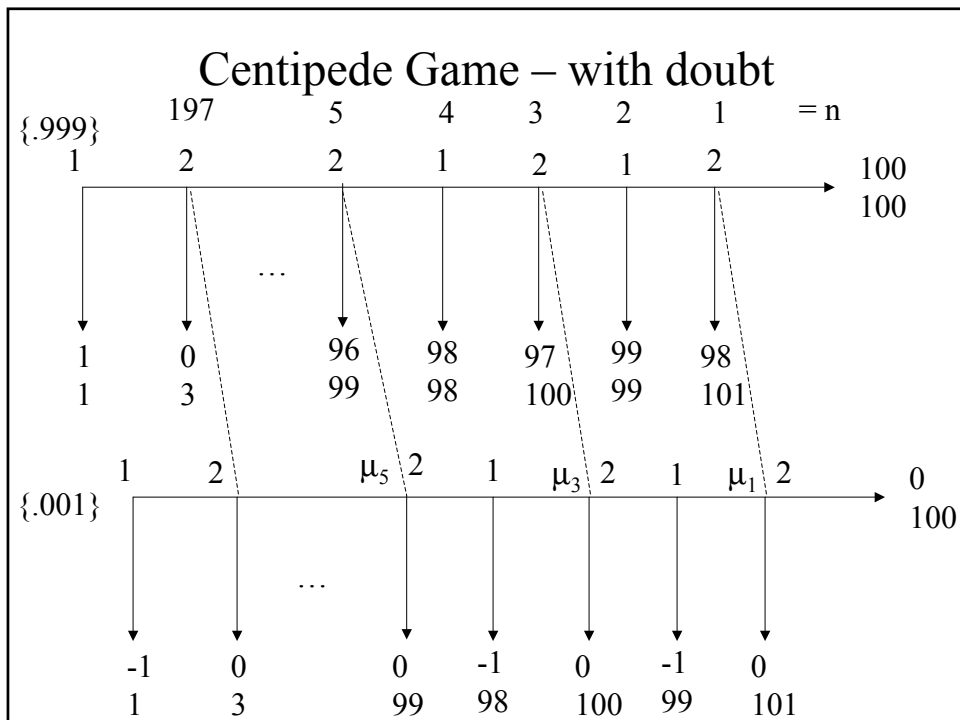


Centipede Game – with doubt



Facts about the Centipede

- Every information set of 2 is reached with positive probability.
- 2 always goes across with positive probability.
- If 2 strictly prefers to go across at n , then
 - 1 must strictly prefer to go across at $n+1$,
 - 2 must strictly prefer to go across at $n+2$,
 - her posterior at n is her prior.
- For any $n > 2$, 1 goes across with positive probability. If 1 goes across w/p 1 at n , then 2's posterior at $n-1$ is her prior.

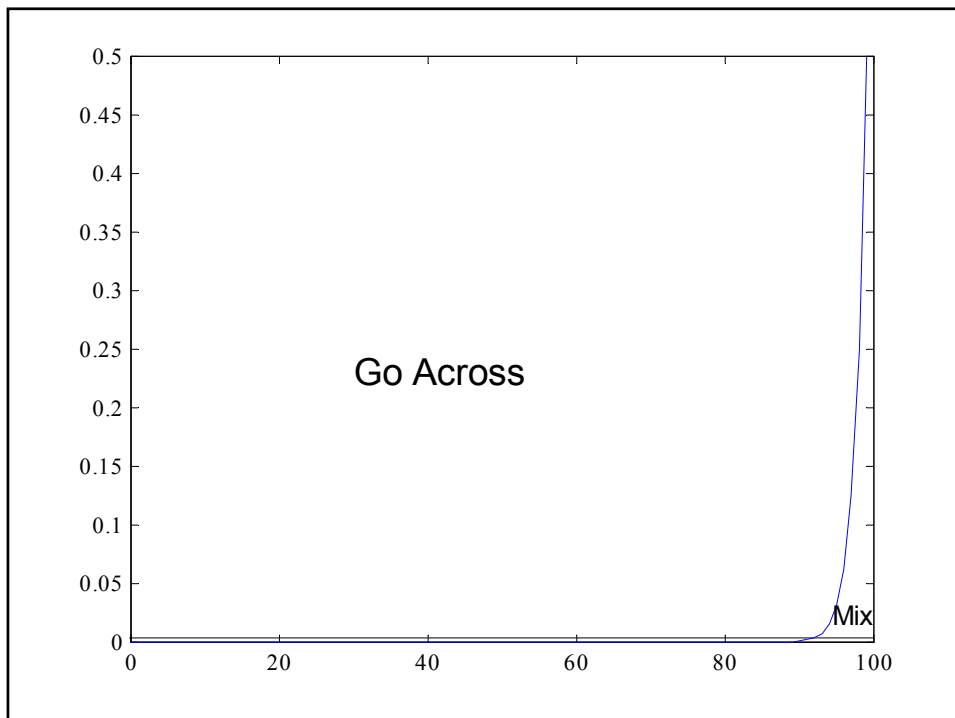


If 2's payoff at any n is x and 2 is mixing,
then

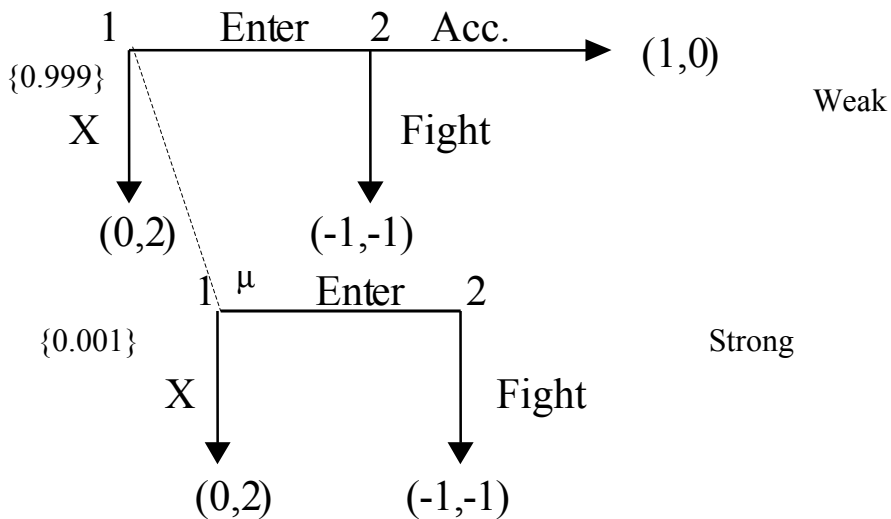
$$\begin{aligned} x &= \mu_n(x+1) + (1-\mu_n)[(x-1)p_n + (1-p_n)(x+1)] \\ &= \mu_n(x+1) + (1-\mu_n)[(x+1) - 2p_n] \\ &= x+1 - 2p_n(1-\mu_n) \\ &\Leftrightarrow (1-\mu_n)p_n = 1/2 \end{aligned}$$

$$\mu_{n-1} = \frac{\mu_n}{\mu_n + (1-\mu_n)(1-p_n)} = \frac{\mu_n}{\mu_n + (1-\mu_n) - p_n(1-\mu_n)} = 2\mu_n$$

$$\mu_n = \frac{\mu_{n-1}}{2}$$



Entry-deterrence with doubt



Entry-deterrence with doubt

- Incumbent is as before.
- Each day there is a new entrant.
- Two types of entrants:
 - W/p $q < 1/2$, tough, $u(\text{enter}, \text{fight}) = 1$;
 - W/p $1-q$, weak, $u(\text{enter}, \text{fight}) = 1$.

Sequential Equilibrium

- Tough entrant (tE) always enters; strong incumbent (sI) always fights.
- If any entrant (E) is accommodated; it becomes common knowledge that incumbent (I) is weak
...
- In the last period,
 - wI accommodates;
 - wE enters iff
- $-1\mu_0 + (1 - \mu_0) \geq 0 \Leftrightarrow \mu_0 \leq 1/2$
- if $\mu_0 < 1/2$, at $n=1$, wI accommodates $\Rightarrow \mu_0 = 1$.
- if $\mu_0 > 1/2$, at $n=1$, wI fights $\Rightarrow \mu_0 = 0.001$.
- $\mu_0 = 1/2$.

- $\mu_0 = \mu_1 / (\mu_1 + \phi_1(1 - \mu_1)) = 1/2$
- $\Leftrightarrow \mu_1 = \phi_1(1 - \mu_1)$
- $\Rightarrow \text{Prob}(\text{fight} | \mu_1) = 2\mu_1$.
- \Rightarrow at $n=1$, wE enters iff $2\mu_1 < 1/2$ i.e., $\mu_1 < 1/4$.
- Similarly, $\mu_1 = 1/4$.
- Proceed as before.