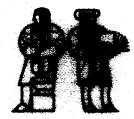
CAMBRIDGE

MASSACHUSETTS 02139

# PROJECT MAC

## MASSACHUSETTS INSTITUTE OF TECHNOLOGY



May 1974

Vera Pless

SYMMETRY CODES AND THEIR INVARIANT SUBCODES

MAC TECHNICAL MEMORANDUM 244

This blank page was inserted to preserve pagination.

Symmetry Codes and Their Invariant Subcodes

#### Abstract

We define and study the invariant subcodes of the symmetry codes in order to be able to determine the algebraic properties of these codes. An infinite family of self-orthogonal rate 1/2 codes over GF(3), called symmetry codes, were constructed in [3]. A (2q + 2, q + 1) symmetry code, denoted by C(q), exists whenever q is an odd prime power  $\equiv$  -1, (mod 3). The group of monomial transformations leaving a symmetry code invariant is denoted by G(q). In this paper we construct two subcodes of C(q) denoted by  $R_{\sigma}(q)$  and  $R_{\mu}(q)$ . Every vector in  $R_{\sigma}(q)$  is invariant under a monomial transformation  $\tau$ in G(q) of odd order s where s divides (q + 1). Also  $R_{_{11}}(q)$  is invariant under  $\tau$  but not vector-wise. The dimensions of  $R_{\sigma}(q)$  and  $R_{\mu}(q)$  are determined and relations between these subcodes are given. An isomorphism is constructed between  $R_{\sigma}(q)$  and a subspace of  $W = V_3 \frac{2q+2}{s}$ . It is shown that the image of  $R_{\sigma}(q)$  is a self-orthogonal subspace of W. The isomorphic images of  $R_{\sigma}(17)$  (under an order 3 monomial) and  $R_{\sigma}(29)$ (under an order 5 monomial) are both demonstrated to be equivalent to the (12, 6) Golay code.

> Dr. Vera Pless Project MAC Massachusetts Institute of Technology 545 Technology Square, Rm. 830 Cambridge, Massachusetts 02139

This empty page was substituted for a blank page in the original document.

## Symmetry Codes and Their Invariant Subcodes by Dr. Vera Pless Project MAC

## I. Introduction.

This paper defines and studies the invariant subcodes of the symmetry codes which were originally defined in [3]. The purpose of this study is the illucidation of properties of these subcodes in such a manner that these properties can be applied in determining characteristics of the symmetry code itself. For example, maximum length vectors in C(17) and C(29) can be determined from known maximum length vectors in the Golay code C(5). The minimum weights are known for the first five symmetry codes. Estimates of the minimum weights of the larger symmetry codes have been obtained by locating a vector of weight 21 in  $R_{_{\rm CT}}(41)$  (under an order 7 monomial) and a vector of weight 27 in  $R_{\sigma}(53)$  (under an order 3 monomial). An (n, k) error correcting code over GF(3) is a k-dimensional subspace of  $V_3^n = V$ . The weight of a vector x, denoted by w(x), is the number of non-zero components it has. Symmetry codes are an infinite family of (2q + 2, q + 1) codes over GF(3) where q is an odd prime power  $\equiv$  -1 (mod 3). Each code is given in terms of a basis [I,  $S_q$ ] where I is the q x q identity matrix and  $S_{q}$  is the matrix described below.

We consider the elements of GF(q) to be ordered in some fixed way, and with this ordering we label the first q + 1 coordinates with the elements of  $GF(q) \cup \{\infty\}$  with  $\infty$  taken as the first coordinate. We label the second q + 1 coordinates by the same sequence of elements of GF(q)  $\bigcup \{\infty\}$  with dashes on them to distinguish them from the first q + 1 coordinate labels. When q = p is a prime, for convenience we use the ordering  $\infty, 0, 1, \ldots$ , p-1 (and hence also  $\infty', 0', 1', \ldots, (p-1)'$  for the right side). By definition, S<sub>q</sub> is the (q + 1) x (q + 1) matrix (s<sub>i'j'</sub>), i, j in GF(q)  $\bigcup \{\infty\}$ , such that s<sub> $\infty', \infty'</sub> = 0 and for i', j' <math>\neq \infty'$ , s<sub>i', $\infty'$ </sub> =  $\chi$  (-1), s<sub> $\infty', i'$ </sub> = 1, and s<sub>i', j'</sub> =  $\chi$ (j-i) where  $\chi(0) = 0, \chi$ (a quadratic residue) = 1, $\chi$  (a non-residue) = -1. We refer to the code generated by [I, S<sub>q</sub>] as C(q).</sub>

As a concrete example we write the basis for C(5) below.

∞ 0	1	2	3	4	<u></u>	0'	1'	2'	3'	_4 <b>'</b>
10	0	0	0	0	0	1	1	1	1	1
0 1	0	0	0	0	1	0	1	-1	-1	1
0 0	1	0	0	0	1	1	0	1	-1	-1
0 0	0	1	0	0	1	-1	1	0	1	-1
0 0	0	0	1	0	1	-1	-1	1	0	1
0 0	0	0	0	1	1	1	-1	-1	1	0

C(5) is a (12, 6) code and it is equivalent to the Golay code [2].

In [4] it was shown that each symmetry code is self orthogonal. The transformations on V which preserve the weights of all vectors are the monomial transformations. A monomial transformation can be viewed as a permutation of the coordinate indices of the vectors in V (the same permutation for each vector) coupled with multiplying some (or none) of the coordinates by minus one. The set of monomial transformations which send all the vectors in C(q) onto vectors in C(q) form a group denoted by G(q). In [4] it was shown that G(q) contains  $PGL_2(q)$ .

In section II of this paper we construct two subcodes of C(q) denoted by  $R_{\sigma}(q)$  and  $R_{\mu}(q)$ . Every vector in  $R_{\sigma}(q)$  is invariant under a monomial transformation  $\tau$  in G(q) of odd order s where s divides q + 1. Also  $R_{\mu}(q)$  is invariant under  $\tau$  but not vector-wise invariant. The dimensions of  $R_{\sigma}(q)$  and  $R_{\mu}(q)$  are determined and relations between these subcodes are given. In section III an isomorphism is constructed between  $R_{\sigma}(q)$  and a subspace of  $W = V_3 \frac{2q+2}{s}$ . It is shown that the image of  $R_{\sigma}(q)$  is a self-orthogonal subspace of W. In section IV the isomorphic images of  $R_{\sigma}(17)$  (o( $\tau$ ) = 3) and  $R_{\sigma}(29)$  (o ( $\tau$ ) = 5), are both demonstrated to be equivalent to the (12, 6) Golay code.

II. In this section we construct two subcodes of C(q), R<sub> $\sigma$ </sub>(q) and R<sub> $\mu$ </sub>(q) with the following properties. Every vector in R<sub> $\sigma$ </sub>(q) is invariant under a monomial transformation  $\tau$  in G(q) where the order of  $\tau$  is an odd number s dividing q + 1. Further, R<sub> $\mu$ </sub>(q) is also invariant under but not vector-wise invariant. The dimensions of R<sub> $\sigma$ </sub> and R<sub> $\mu$ </sub> are determined, and relations between them are given.

In [4] it was shown that the mapping sending a monomial transformation  $\tau$  in G(q) onto the permutation  $\overline{\tau}$  it induces on the coordinate indices is a homomorphism of a subgroup of G(q) onto PGL<sub>2</sub>(q) whose kernel has order 2. For the rest of this paper  $\tau$  denotes a monomial transformation in G(q) of odd order s where s divides (q + 1) such that  $\overline{\tau}$  is in PGL<sub>2</sub>(q) and the order of  $\tau$  equals the order of  $\overline{\tau}$ .

Lemma 1. If s is an odd number dividing (q + 1), then there exists a transformation  $\overline{\tau}$  in G(q) or order s. Further  $\overline{\tau}$  is in PGL<sub>2</sub>(q).

Proof: By [1] it is known that  $PGL_2(q)$  contains a cyclic subgroup of order  $(\frac{q+1}{2})$ . Hence this subgroup contains an element  $\overline{\tau}$  of order s when s is any odd number dividing (q + 1). The monomial  $\tau$  in G(q) which maps into  $\overline{\tau}$  by the homomorphism described above is either of order s or 2s. If it is of order s we are finished. If  $\tau$  is of order 2s then  $\tau^2$ is of order s,  $\overline{\tau}^2$  is also of order s (since s is odd),  $\overline{\tau}^2$  is in  $PGL_q(q)$ and the lemma is demonstrated.

The subcodes R  $_{\sigma}(q)$  and R  $_{\mu}(q)$  are the ranges of two linear transformations  $\sigma$  and  $_{\mu}$  defined for x in C(q) as follows.

 $x\sigma = x + x\tau + \dots + x\tau^{s-1}$  $x\mu = x - x\tau$ 

Even though  $\sigma$  and  $\mu$  are linear transformations, they are not monomial transformations; they are useful in obtaining information about  $\tau$ . Let  $K_{\sigma}(q)$  denote the kernel of  $\sigma$  and  $K_{\mu}(q)$  the kernel of  $\mu$ .

Theorem 1.  $R_{\sigma}(q)$ ,  $R_{\mu}(q)$ ,  $K_{\sigma}(q)$ ,  $K_{\mu}(q)$  are subcodes of C(q) such that

1)  $R_{\sigma}(q)$  is contained in  $K_{\mu}(q)$  and  $R_{\mu}(q)$  is contained in  $K_{\sigma}(q)$ , and 2)  $\tau$  leaves  $R_{\mu}(q)$  invariant and  $\tau$  leaves every vector in  $R_{\sigma}(q)$ invariant.

Proof: It is clear that  $R_{\sigma}(q)$ ,  $R_{\mu}(q)$ ,  $K_{\sigma}(q)$  are subcodes since they are vector subspaces contained in C(q). If  $x\sigma$  is in  $R_{\sigma}(q)$  then  $(x\sigma)\mu =$  $(x + x\tau + ... + x\tau^{s-1})_{\mu} = (x + x\tau + ... x\tau^{s-1}) - (x\tau + x\tau^2 + ... + x\tau^{s-1} + x) = 0$  so that  $R_{\sigma}(q)$  is contained in  $K_{\mu}(q)$ . Similarly  $R_{\mu}(q)$  is contained in  $K_{\sigma}(q)$ . If  $x\sigma$  is in  $R_{\sigma}(q)$ , then  $(x\sigma)\tau = (x + x\tau + ... + x\tau^{s-1})\tau =$  $x\tau + x\tau^2 + ... x\tau^{s-1} + x = x\sigma$  and we see that  $\tau$  leaves every vector in  $R_{\sigma}(q)$  invariant. Since  $(x\mu)\tau = x\tau - x\tau^2$ ,  $\tau$  leaves  $R_{\mu}(q)$  invariant and the theorem is proved. Remark: When s is divisible by 3,  $R_{\sigma}(q)$  is contained in  $K_{\sigma}(q)$ . Proof: If y is in  $R_{\sigma}(q)$ ,  $y = x\sigma = x + x\tau + \dots + x\tau^{s-1}$ . Hence  $y\sigma = (x + x\tau + \dots x\tau^{s-1})\sigma = sy \equiv 0 \pmod{3}$ .

Lemma 2.  $\overline{\tau}$  is a product of disjoint cycles of length s. Further, if ( $i_1$ , ...,  $i_s$ ) is such an s-cycle for the left coordinate indices of V, then ( $i_1$ ', ...,  $i_s$ ') is such an s-coch cle for the right coordinate indices of V.

Proof: By their construction [4] the transformations in  $PGL_2(q)$  act on the left coordinate indices (and simultaneously on the right coordinate indices) as transformations on the projective line. Since s is an odd number which divides q + 1,  $\overline{\tau}$  is either completely a product of disjoint cycles of length s or a product of disjoint cycles of length s with ks fixed points. But a projective transformation with three fixed points is the identity. Hence  $\overline{\tau}$  can have at most two fixed points on each side of coordinate indices. Since s divides q + 1, the number of left coordinate indices (and the number of right coordinate indices), this is only possible for k = 1 and s = 2. The lemma follows from the fact that s is an odd number.

We let J be a set of left coordinate indices with the property that J contains exactly one index from each of these s cycles. Note that  $|J| = \frac{(q+1)}{s}$ .

In order to determine the dimension of  $R_{\sigma}(q)$  and  $R_{\mu}(q)$  we introduce the following terminology. We let the vectors in the basis [I,  $S_q$ ] be denoted by ( $e_i$ ,  $c(e_i)$ ) where  $e_i$  is the  $i\frac{th}{t}$  row of I and  $c(e_i)$  is the  $i\frac{th}{t}$ row of  $S_q$ . Theorem 2. dim  $R_{\sigma}(q) = \frac{(q+1)}{s}$  and dim  $R_{\mu}(q) = \frac{(q+1)(s-1)}{s}$ . Proof: Consider the set of  $\frac{(q+1)}{s}$  vectors  $\{(e_j + e_j\tau + \ldots + e_j\tau^{s-1}, c(e_j) + c(e_j)\tau + \ldots + c(e_j)\tau^{s-1})\}$  for jeJ. Since the order of  $\tau$  equals the order of  $\overline{\tau}$ ,  $e_j \neq \frac{t}{r} = e_j \tau^i$ ,  $1 \leq i \leq s - 1$ , so that  $(e_j + e_j\tau + \ldots + e_j\tau^{s-1}) \neq 0$  for each j  $\epsilon$  J. Hence by the definition of J, these vectors are linearly independent. Clearly they span  $R_{\sigma}(q)$ , and it thus follows that dim  $R_{\sigma}(q) = |J| = \frac{q+1}{s}$ . Similarly  $\{(e_j\tau^k - e_j\tau^{k+1}), (c(e_j)\tau^k - c(e_j)\tau^{k+1})\}$  for jeJ,  $k = 0, \ldots, s - 2$  is a basis of  $R_{\mu}(q)$ . Hence dim  $R_{\mu}(q) = \frac{(q+1)(s-1)}{s}$ .

Remark: When  $\tau$  has even order ( $\neq 2$ ) which divides  $(\frac{q+1}{2})$ , all the results of this paper hold when the order of  $\tau$  equals the order of  $\overline{\tau}$ . When the order of  $\tau$  equals twice the order of  $\overline{\tau}$ , then it is possible that Theorem 2 does not hold since the basis vectors described above can be zero.

Corollary 1.  $R_{\sigma}(q) = K_{\mu}(q)$  and  $R_{\mu}(q) = K_{\sigma}(q)$ . Proof: By Theorem 1,  $R_{\mu}(q)$  is contained in  $K_{\sigma}(q)$  and  $R_{\sigma}(q)$  is contained in  $K_{\mu}(q)$ . In general, dim  $R_{\mu}(q) + \dim K_{\mu}(q) = q + 1 = \dim K_{\sigma}(q) + \dim R_{\sigma}(q)$ . By Theorem 2, dim  $R_{\sigma}(q) = \frac{(q+1)}{s}$  and dim  $R_{\mu}(q) = \frac{(q+1)(s-1)}{s}$ . Hence dim  $R_{\mu}(q) = \dim K_{\sigma}(q)$  and dim  $R_{\sigma}(q) = \dim K_{\mu}(q)$  and the corollary is demonstrated.

Note that since 3 divides (q + 1) for every  $q \equiv -1 \pmod{3}$ , every symmetry code has a monomial transformation of order 3 leaving it invariant.

III. The isomorphic image of  $R_{\sigma}$ .

In this section we construct a linear transformation  $\varphi$  from V onto W = V<sub>3</sub>  $\frac{2q+2}{s}$  where s is again an odd number dividing q + 1 with the following properties. The dimension of  $\varphi(R_{\sigma})$  equals the dimension of  $R_{\sigma}$ , the weight of  $\varphi(x)$  for x in  $R_{\sigma}$  is the weight of x divided by s, and  $\varphi(R_{\sigma})$  is a self-orthogonal subspace of W.

In order to do this we let J be as in section II, and let J' be the elements in J with dashes on them. Note that  $J \cup J'$  contains  $\frac{2(q+1)}{s}$  elements. We consider the elements in J to have the same ordering they had in  $GF(q) \cup \{\infty\}$ . With this ordering we label the left half of the coordinate indices in W with the elements from J, and the right half with the elements from J. We denote the unit vectors in W by  $\overline{e_j}$ , j in J and  $\overline{e_j'}$ , j' in J'.

Lemma 3. If  $x\tau = x$ , then the components of x on a cycle of  $\overline{\tau}$  are either all zero or all non-zero. Further, if  $x\tau = x$  and  $y\tau = y$ , then on the cycles of  $\overline{\tau}$  on which the components of both x and y are non-zero, the components of x equal plus or minus the components of y. Proof: Let  $(i_1, \dots, i_s)$  be the coordinate indices of a cycle of  $\overline{\tau}$ . Let  $x_{i_j}$  be the  $i_j \frac{th}{t}$  component of x. If  $x\tau = x$ , then all the components of x on this cycle are determined by  $x_{i_1}$  and  $\tau$ . If  $y\tau = y$  also, then the components of x on this cycle equal the components of y on this cycle of  $x_{i_1} =$  $y_{i_1}$ . If  $x_{i_1} = -y_{i_1}$  the components of x on this cycle are the negatives of the components of y. Since these are the only possibilities, the lemma is proved.

Theorem 3. There is a linear transformation  $\varphi$  from V onto W = V<sub>3</sub> such that 1) dim  $\varphi(R_{\sigma}(q)) = \dim R_{\sigma}(q) = \frac{(q+1)}{s}$ , and 2)  $w(\varphi(x)) = \frac{w(x)}{s}$ .

-7-

Proof: We let  $e_i$  and  $e_i'$ ,  $i \in GF(q) \cup \{\infty\}$  denote the unit vectors in V. We define  $\varphi$  on these unit vectors as follows.

If 
$$j \in J$$
,  $\varphi(e_j) = \overline{e_j}$ . If  $i \notin J$ ,  $\varphi(e_i) = 0$ .  
If  $j' \in J'$ ,  $\varphi(e_j') = \overline{e_j'}$ . If  $i' \notin J'$ ,  $\varphi(e_i') = 0$ .

Define  $\phi$  on the rest of V linearly. Clearly  $\phi$  is a linear transformation from V onto W.

Recall that  $\{(e_j + e_j \tau + ... + e_j \tau^{s-1}, c(e_j) + c(e_j) \tau + ... + c(e_j) \tau^{s-1})\}$ , jeJ is a basis of  $R_{\sigma}(q)$ . Since  $\varphi$  maps these vectors onto linearly independent vectors, dim  $\varphi(R_{\sigma}(q)) = \dim R_{\sigma}(q) = \frac{(q+1)}{s}$  by Theorem 2.

Theorem 1 tells us that  $x\tau = x$  for all x in  $R_{\sigma}(q)$ . By Lemma 3 we know that the components of x on a cycle of  $\overline{\tau}$  are either all zero or all non-zero. Since  $\varphi$  projects on precisely one component from each s-cycle of  $\overline{\tau}$ ,  $w(\varphi(x)) = \frac{w(x)}{s}$ .

It was proven in [4] that C(q) is a self-orthogonal subspace of V so that  $R_{\sigma}(q)$  is certainly a self-orthogonal subspace of V. Even though  $\varphi$ does not preserve the property of self-orthogonality, we can prove that  $\varphi(R_{\sigma}(q))$  is a self-orthogonal subspace of W.

Theorem 4.  $\varphi(\mathsf{R}_{\sigma}(\mathsf{q}))$  is a self-orthogonal subspace of W. Proof: Let x and y be vectors in W such that  $x = (\alpha_1, \dots, \alpha_{\underline{2q+2}})$  and  $\underline{s}$ 

 $y = (\beta_1, \dots, \beta_{\frac{2q+2}{s}})$ . Then the inner product of x and y, denoted by (x,y), is  $\begin{pmatrix} \frac{2q+2}{s} \\ \sum_{i=1}^{s} \alpha_i \beta_i \end{pmatrix}$  (mod 3). As is usual, x and y are orthogonal to each other if (x,y) = 0. In order to prove Theorem 4 we need to show that (x,y) = 0for all x,y in  $\varphi(R_{\sigma}(q))$  (x can also equal y). In order to prove this, we introduce the inner product of x and y over the integers, denoted by  $\frac{2q+2}{s}$ [x,y], where [x,y] equals  $\sum_{i=1}^{n} \alpha_i \beta_i$  by definition. We define [x,y] in i=1

a similar fashion for x and y in V.

The proof of Theorem 4 is divided into two cases. The first case is 3 does not divide s. If x and y are in  $R_{\sigma}(q)$ , then  $x = x_1 + x_1\tau + \ldots + x_1\tau^{s-1}$  and  $y = y_1 + y_1\tau + \ldots + y_1\tau^{s-1}$  for some  $x_1$  and  $y_1$  in C(q). By Lemma 3, all the elements in  $R_{\sigma}(q)$  which are not zero on a particular cycle of  $\overline{\tau}$  have the same or opposite components on that cycle. Hence [x,y] = rswhere r is the number of s-cycles of  $\overline{\tau}$  (in both the left and right coordinates) in which both x and y have non-zero components. Since (x,y) = 0, 3 divides rs, but by assumption 3 does not divide s so that 3 divides r. By the definition of  $\varphi$ ,  $[\varphi(x), \varphi(y)] = r$  so that  $(\varphi(x), \varphi(y)) = 0$  for all x,y in  $R_{\sigma}(q)$ . Hence  $\varphi(R_{\sigma}(q))$  is self-orthogonal in this situation. We now consider the case that s = 3j, i.e.,  $\tau^{3j} = 1$ . We let x and y be in  $R_{\sigma}(q)$ , and we have  $x = x_1 + x_1^{\tau} + \ldots + x_1\tau^{3j-1}$ ,  $y = y_1 + y_1\tau + \ldots + y_1\tau^{3j-1}$  for  $x_1, y_1$  in C(q). Then

$$[x,y] = \sum_{i=0}^{3j-1} [x_1, y_1\tau^{i}] + \sum_{i=0}^{3j-1} [x_1\tau, y_1\tau^{i}] + \dots + \sum_{i=0}^{3j-1} [x_1\tau^{3j-1}, y_1\tau^{i}]$$

$$= \sum_{i=0}^{3j-1} [x_1\tau^{i}, y_1\tau^{i}] + \sum_{i=0}^{3j-1} [x_1\tau^{i}, y_1\tau^{i+1}] + \dots + \sum_{i=0}^{3j-1} [x_1\tau^{i}, y, \tau^{i+3j-1}]$$

by rearranging terms. Now  $[u,v] = [u\tau^{i}, v\tau^{i}]$  for all u and v in V

-9-

since  $\tau^{i}$  is a monomial transformation over GF(3). Hence  $[x,y] = 3j[x_{1}, y_{1}] + 3j[x_{1}, y_{1}\tau] + \ldots + 3j[x_{1}, y_{1}\tau^{3j-1}]$ . Since  $x_{1}$  and  $y_{1}\tau^{i}$  (i=0,...,3j-1) are all in C(q) which is self-orthogonal, each  $[x_{1}, y_{1}\tau^{i}]$  is divisible by 3 so that [x,y] = 9r for some r. Each cycle of  $\overline{\tau}$  is a 3j-cycle, and by the definition of  $\varphi$ ,  $\varphi$  projects onto one coordinate from each 3j-cycle so that  $[\varphi(x), \varphi(y)] = 3r$ . Hence  $(\varphi(x), \varphi(y)) = 0$ , and  $\varphi(R_{\sigma}(q))$  is a self-orthogonal subspace of W for this case also.

IV. Invariant subcodes of C(17) and C(29) are isomorphic to the Golay code.

In this section we apply these ideas to C(17) and C(29). The  $\tau$  for C(17) has order 3 and the  $\tau$  for C(29) has order 5. We describe these two monomial transformations explicitly, and exhibit bases for  $R_{\sigma}(17)$  and  $\cdot \phi(R_{\sigma}(17))$ .

In order to exhibit these monomial transformations we introduce the following convention. We let  $\overline{\chi(i)}$  times a column index mean that we multiply the column by  $\chi(i)$  where  $\chi(i) = 1$  for i a quadratic residue, and  $\chi(i) = -1$  for i a non-residue. This convention is used in order to avoid confusion with negatives in GF(17).

We can represent  $\tau$  as a monomial transformation on the columns of V as follows.

$$\tau(\infty) = 0, \quad \tau(16) = \infty; \quad \tau(i) = \overline{\chi(i+1)} \quad \left(\frac{16}{i+1}\right), \quad i \neq \infty, \quad 16;$$
  
$$\tau(\infty') = 0', \quad \tau(16') = \infty'; \quad \tau(i') = \overline{\chi(i'+1)} \quad \left(\frac{16}{i+1}\right), \quad i' \neq \infty', \quad 16'.$$

The generators of the subgroup of G(17) which is isomorphic to  $PGL_2(17)$  are given in [4, p. 131]. It is easy to verify that  $\tau$  is a product of two

of these generators so that  $\tau$  is in G(17). A straightforward check shows that  $\tau$  has order 3. If we rearrange the columns of V to correspond to the cycles of  $\overline{\tau}$ , the following is a basis of R<sub>0</sub>(17).

<u>∞ 0 16</u>	1 8 15	2 11 7	3 4 10	5 14 9	6 12 13	<u>∞' 0' 16'</u>	1' 8' 15'	2' 11' 7'	3' 4' 10'	5' 14' 9'	6' 12' 13'
111				1		-1 -1 -1		1 -1 1	1 1 -1	-1 1 1	-1 1 1
	1 1 1						1 1 1	1 -1 1	-1 -1 1	-1 1 1	1 -1 -1
		1 -1 1				1 1 1	1 1 1	1 -1 1	1 1 -1	1 -1 -1	l
			11-1			1 1 1	-1 -1 -1	1 -1 1	-1 -1 1		-1 1 1
				1 -1 -1		-1 -1 -1	-1 -1 -1	1 -1 1		1 - 1 - 1	1 -1 -1
					1 -1 -1	-1 -1 -1	1 1 1		-1 -1 1	1 -1 -1	-1 1 1

From this we get the following basis for  $\varphi(R_{\sigma}(17))$  by choosing  $J = \{\infty, 1, 2, 3, 5, 6\}$ .

8	1	2	3	5	6	<b>∞'</b>	1'	2'	3'	5'	6'
1						-1		1	1	-1	-1
	1						1	1	-1	-1	1
		1				1	1	1	1	1	
			1			1	-1	1	-1		-1
				1		-1	-1	1		1	1
					1	-1	1		-1	1	-1

It is known [4] that the minimum weight of C(17) is 18, so that the minimum weight of  $\varphi(R_{\sigma}(17))$  is 6. It follows from the theorem in [2] that  $\varphi(R_{\sigma}(17))$  is equivalent to the Golay (12, 6) code over GF(3).

A monomial transformation  $\tau$  of order 5 in G(29) is given by the following.

$$\tau(\infty) = 0, \ \tau(24) = \infty; \ \tau(i) = \overline{\chi(i+5)} \left(\frac{28}{i+5}\right), \ i \neq \infty, \ 24,$$
  
$$\tau(\infty') = 0', \ \tau(24') = \infty'; \ \tau(i') = \overline{\chi(i'+5)} \left(\frac{28}{i'+5}\right), \ i' \neq \infty', \ 24'.$$
  
As in the previous case it can be verified that  $\tau$  is a product of

generators of the subgroup of G(29) which is isomorphic to PGL<sub>2</sub>(29). Given  $\tau$ , a basis of  $R_{\sigma}(29)$  can be computed similar to the basis of  $R_{\sigma}(17)$ . The minimum weight in C(29) is 18 and since the weight of every vector in  $R_{\sigma}(29)$  is divisible by 5, the minimum weight of  $R_{\sigma}(29)$  must be at least 30. It is empethy 80 since the basis vectors have weight 30. Hence the minimum weight of  $\sigma(29)$  is 6. It then follows as above that  $\sigma(29)$  is equivalent to the delay Gode.

I wish to thank Jean-Marie Goethals for pointing out to me that the results of this paper are applicable to a wider class of monomials

as an this for a construction and takes the provide the

and the second of the second second

he should be a

e e le letter e le men de le recidente de le marca A

and the state of the

and and the formation and the standard standard and the standard standard and and

than I originally stated.

## Bibliography

- 1. Dickson, L. E. (1901, 1958) "Linear Groups with an Exposition of the Galois Field Theory", reprinted by Dover Publications, New York.
- 2. V. Pless, "On the uniqueness of the Golay codes", J. of Combinatorial Theory, 5 (1968), 215-228.
- V. Pless, "On a new family of symmetry codes and related new fivedesigns", <u>Bulletin of the American Mathematical Society</u>, Vol. 75, No. 6 (1969), 1339-1342.
- 4. V. Pless, "Symmetry codes over GF(3) and new five-designs", <u>J. of</u> <u>Combinatorial Theory</u>, 12 (1972), 119-142.

BIBLIOGRAPHIC DATA SHEET 1. Report No. MAC TM-44	2. 3. Recipient's Accession No.
4. Title and Subtitle	5. Report Date : Issued
Symmetry Codes and their Invariant Subcodes	s <u>May 1974</u>
	0.
7. Author(s) Vera Pless	<ol> <li>Performing Organization Rept.</li> <li>No. MAC TM- 44</li> </ol>
9. Performing Organization Name and Address	10. Project/Task/Work Unit No.
PROJECT MAC; MASSACHUSETTS INSTITUTE OF TEC	CHNOLOGY: 11. Contract/Grant No.
545 Technology Square, Cambridge, Massachus	
	N00014-70-A-0362-0006
12. Sponsoring Organization Name and Address Office of Naval Research	<b>13.</b> Type of Report & Period Covered <b>Interim</b>
Department of the Navy	Scientific Report
Information Systems Program	14.
Arlington, Va 22217	
15. Supplementary Notes	
16. Abstracts : The paper defines and studies th	he invariant subcodes, $R_{\sigma}(q)$ and $R_{\mu}(q)$ , of
the symmetry code $C(q)$ in order to be able these codes. Every vector in $R(q)$ is inva	to determine the algebraic properties of ariant under a monomial transformation $ au$ , of
odd order dividing $(q + 1)$ , in the group of	
but not vector-wise. The dimensions of R	
-	3
between these subcodes are given. Also R $\sigma$	(q) is shown to be isomorphic to a self-
orthogonal subspace of $V_3$ s. The isome	orphic images of $ m R_{\sigma}(17)$ and $ m R_{\sigma}(29)$ are both
demonstrated to be equivalent to the (12,6)	) Golay code.
17. Key Words and Document Analysis. 170. Descriptors	
•	
17b. Identifiers/Open-Ended Terms	
17c. COSATI Field/Group	
18. Availability Statement	19. Security Class (This 21. No. of Pages
Approved for Public Release;	Report) UNCLASSIFIED 16
Distribution Unlimited	20. Security Class (This 22. Price
	Page UNCLASSIFIED
FORM NTIS-35 (REV. 3-72) THIS FORM MAY B	USCOMM-DC 14952-P72

ι.

-

•

.

-