
Learning—Adjustment with persistent noise

14.126 Game Theory
Mihai Manea
Muhamet Yildiz

Main idea

- There will always be small but positive probability of mutation.
 - Then, some of the strict Nash equilibria will **not** be “stochastically stable.”
-

General Procedure

Stochastic Adjustment

1. Consider a game.
2. Specify a state space Θ , e.g., the # of players playing a strategy.

1.

	A	B
A	2,2	0,0
B	0,0	1,1

2. $\Theta = \{AA, AB, BA, BB\}$

Stochastic Adjustment, continued

3. Specify an adjustment dynamics, e.g., best-response dynamics, with a transition matrix P , where

$$P_{\theta,\xi} = \Pr(\theta \text{ at } t+1 | \xi \text{ at } t)$$

ϕ = a probability distribution, a column vector.

3.

		AA	AB	BA	BB	
$P =$	1	0	0	0	0	AA
	0	0	1	0	0	AB
	0	1	0	0	0	BA
	0	0	0	1	1	BB

Stochastic Adjustment, continued

4. Introduce a small noise: Consider P^ε , continuous in ε and $P^\varepsilon \rightarrow P$ as $\varepsilon \rightarrow 0$.

Make sure that there exist a unique ϕ_ε^* s.t.

$$\phi_\varepsilon^* = P^\varepsilon \phi_\varepsilon^*.$$

4. AA AB BA BB

$P^\varepsilon =$	$(1-\varepsilon)^2$	$(1-\varepsilon)\varepsilon$	$(1-\varepsilon)\varepsilon$	ε^2
	$(1-\varepsilon)\varepsilon$	ε^2	$(1-\varepsilon)^2$	$(1-\varepsilon)\varepsilon$
	$(1-\varepsilon)\varepsilon$	$(1-\varepsilon)^2$	ε^2	$(1-\varepsilon)\varepsilon$
	ε^2	$(1-\varepsilon)\varepsilon$	$(1-\varepsilon)\varepsilon$	$(1-\varepsilon)^2$

$$\phi_\varepsilon^* = (1/4, 1/4, 1/4, 1/4)^\top.$$

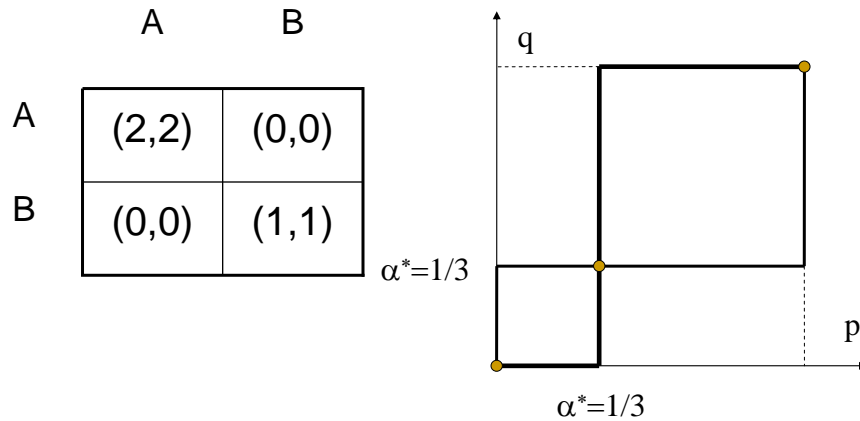
Stochastic Adjustment, continued

5. Verify that $\lim_{\varepsilon \rightarrow 0} \phi_\varepsilon^* = \phi^*$ exists; compute ϕ^* .
(By continuity $\phi^* = P\phi^*$.)
6. Check that ϕ^* is a point mass, i.e.,
$$\phi^*(\theta^*) = 1$$
for some θ^* .

The strategy profile at θ^* is called *stochastically stable equilibrium*.

Kandoori, Mailath & Rob

Coordination game



Adjustment Process

- N = population size.
- θ_t = # of players who play A at t .
- $u_A(\theta_t) = \theta_t/N u(A,A) + (N - \theta_t)/N u(A,B)$
- $\theta_{t+1} = P(\theta_t)$, where
 - $P(\theta_t) > \theta_t \Leftrightarrow u_A(\theta_t) > u_B(\theta_t)$ &
 - $P(\theta_t) = \theta_t \Leftrightarrow u_A(\theta_t) = u_B(\theta_t)$.

- Example:

$$P(\theta_t) = BR(\theta_t) = \begin{cases} N & \text{if } u_A(\theta_t) > u_B(\theta_t) \\ \theta_t & \text{if } u_A(\theta_t) = u_B(\theta_t) \\ 0 & \text{if } u_A(\theta_t) < u_B(\theta_t) \end{cases}$$

Noise

- Independently, each agent with probability 2ε mutates, and plays either of the strategies with equal probabilities.

$$P^\varepsilon = \begin{bmatrix} (1-\varepsilon)^N & (1-\varepsilon)^N & \dots & (1-\varepsilon)^N & \varepsilon^N & \dots & \varepsilon^N \\ N(1-\varepsilon)^{N-1}\varepsilon & N(1-\varepsilon)^{N-1}\varepsilon & \dots & N(1-\varepsilon)^{N-1}\varepsilon & N(1-\varepsilon)\varepsilon^{N-1} & \dots & N(1-\varepsilon)\varepsilon^{N-1} \\ \vdots & \vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\ N(1-\varepsilon)\varepsilon^{N-1} & N(1-\varepsilon)\varepsilon^{N-1} & \dots & N(1-\varepsilon)\varepsilon^{N-1} & N(1-\varepsilon)^{N-1}\varepsilon & \dots & N(1-\varepsilon)^{N-1}\varepsilon \\ \varepsilon^N & \varepsilon^N & \dots & \varepsilon^N & (1-\varepsilon)^N & \dots & (1-\varepsilon)^N \end{bmatrix}$$

- $\phi^*(\varepsilon) =$ invariant distribution for P^ε .

P^ε

$$\begin{array}{ccc|ccc} (1-\varepsilon)^N & \dots & (1-\varepsilon)^N & \varepsilon^N & \dots & \varepsilon^N \\ N(1-\varepsilon)^{N-1}\varepsilon & \dots & N(1-\varepsilon)^{N-1}\varepsilon & N(1-\varepsilon)\varepsilon^{N-1} & \dots & N(1-\varepsilon)\varepsilon^{N-1} \\ \vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\ \binom{N}{N^*-1}(1-\varepsilon)^{N-N^*+1}\varepsilon^{N^*-1} & \dots & \binom{N}{N^*-1}(1-\varepsilon)^{N-N^*+1}\varepsilon^{N^*-1} & \binom{N}{N^*-1}(1-\varepsilon)^{N^*-1}\varepsilon^{N-N^*+1} & \dots & \binom{N}{N^*-1}(1-\varepsilon)^{N^*-1}\varepsilon^{N-N^*+1} \\ \hline \binom{N}{N^*}(1-\varepsilon)^{N-N^*}\varepsilon^{N^*} & \dots & \binom{N}{N^*}(1-\varepsilon)^{N-N^*}\varepsilon^{N^*} & \binom{N}{N^*}(1-\varepsilon)^{N^*}\varepsilon^{N-N^*} & \dots & \binom{N}{N^*}(1-\varepsilon)^{N^*}\varepsilon^{N-N^*} \\ \vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\ \varepsilon^N & \dots & \varepsilon^N & (1-\varepsilon)^N & \dots & (1-\varepsilon)^N \end{array}$$

$$N^* = \lceil a^*N \rceil < N/2.$$

Invariant Distribution

- $N^* = \lceil \alpha^* N \rceil < N/2$.
- $D_A = \{\theta | \theta \geq N^*\}$; $D_B = \{\theta | \theta < N^*\}$;
- $q_{AB} = \Pr(\theta_{t+1} \in D_A | \theta_t \in D_B)$;
- $q_{BA} = \Pr(\theta_{t+1} \in D_B | \theta_t \in D_A)$;
- $$p_A(\varepsilon) = \sum_{\theta \in D_A} \varphi_\varepsilon^*(\theta)$$
- $p_B(\varepsilon) = 1 - p_A(\varepsilon)$.

Invariant distribution, continued

- $$\begin{bmatrix} p_A(\varepsilon) \\ p_B(\varepsilon) \end{bmatrix} = \begin{bmatrix} 1 - q_{BA} & q_{AB} \\ q_{BA} & 1 - q_{AB} \end{bmatrix} \begin{bmatrix} p_A(\varepsilon) \\ p_B(\varepsilon) \end{bmatrix}$$
- $$\begin{bmatrix} -q_{BA} & q_{AB} \\ q_{BA} & -q_{AB} \end{bmatrix} \begin{bmatrix} p_A(\varepsilon) \\ p_B(\varepsilon) \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$
- $$\frac{p_A(\varepsilon)}{p_B(\varepsilon)} = \frac{q_{AB}}{q_{BA}}$$

Invariant distribution, continued

$$\begin{aligned}
 \frac{p_A(\varepsilon)}{p_B(\varepsilon)} &= \frac{q_{AB}}{q_{BA}} \\
 &= \frac{\binom{N}{N^*} \varepsilon^{N^*} (1-\varepsilon)^{N-N^*} + \binom{N}{N^*+1} \varepsilon^{N^*+1} (1-\varepsilon)^{N-N^*-1} + \dots}{\binom{N}{N^*-1} \varepsilon^{N-N^*+1} (1-\varepsilon)^{N^*-1} + \binom{N}{N^*-2} \varepsilon^{N-N^*+2} (1-\varepsilon)^{N^*-2} + \dots} \\
 &= \frac{\binom{N}{N^*} \varepsilon^{N^*} (1-\varepsilon)^{N-N^*} + o(\varepsilon^{N^*})}{\binom{N}{N^*-1} \varepsilon^{N-N^*+1} (1-\varepsilon)^{N^*-1} + o(\varepsilon^{N-N^*+1})} \cong \frac{\binom{N}{N^*}}{\binom{N}{N^*-1}} \frac{1}{\varepsilon^{N-2N^*+1}} \rightarrow \infty.
 \end{aligned}$$

Proposition

If N is large enough so that $N^* < N/2$, then limit φ^* of invariant distributions puts a point mass on $\theta = N$, corresponding to all players playing A.

Replicator dynamics & Evolutionary stability

14.126 Game Theory
Muhamet Yildiz

Road Map

1. Evolutionarily stable strategies
 2. Replicator dynamics
-

Notation

- $G = (S, A)$ a symmetric, 2-player game where
 - S is the strategy space;
 - $A_{i,j} = u_1(s_i, s_j) = u_2(s_j, s_i)$.
- x, y are mixed strategies; $u(x, y) = x^T A y$;
 $u(s, y)$.
- $ax + (1-a)y$.
- $u(ax + (1-a)y, z) = au(x, z) + (1-a)u(y, z)$
- $u(x, ay + (1-a)z) = au(x, y) + (1-a)u(x, z)$

ESS

Definition: A (mixed) strategy x is said to be *evolutionarily stable* iff, given any $y \neq x$, there exists $\varepsilon_y > 0$ s.t.

$$u(x, (1-\varepsilon)x + \varepsilon y) > u(y, (1-\varepsilon)x + \varepsilon y)$$

for each ε in $(0, \varepsilon_y]$.

- Each player is endowed with a (mixed) strategy.
- Assumes that population is a state
- Asks whether a strategy (state) is robust to evolutionary pressures.
- Disregards effects on future actions.

Alternative Definition





Fact: x is evolutionarily stable iff, $\forall y \neq x$,

1. $u(x,x) \geq u(y,x)$, and
2. $u(x,x) = u(y,x) \Rightarrow u(x,y) > u(y,y)$.

Proof: Define

$$\begin{aligned} F(\varepsilon, y) &= u(x, (1-\varepsilon)x + \varepsilon y) - u(y, (1-\varepsilon)x + \varepsilon y) \\ &= u(x-y, (1-\varepsilon)x + \varepsilon y) \\ &= (1-\varepsilon) u(x-y, x) + \varepsilon u(x-y, y). \end{aligned}$$

Hawk-Dove game

		
	(1-c, 1-c)	(2, 0)
	(0, 2)	(1, 1)

1. $c < 1$
2. $c > 1$

ESS-NE

- If x is an ESS, then (x,x) is a Nash equilibrium.
- In fact, (x,x) is a proper equilibrium.
- If (x,x) is a strict Nash equilibrium, then x is ESS.

Rock-Scissors-Paper

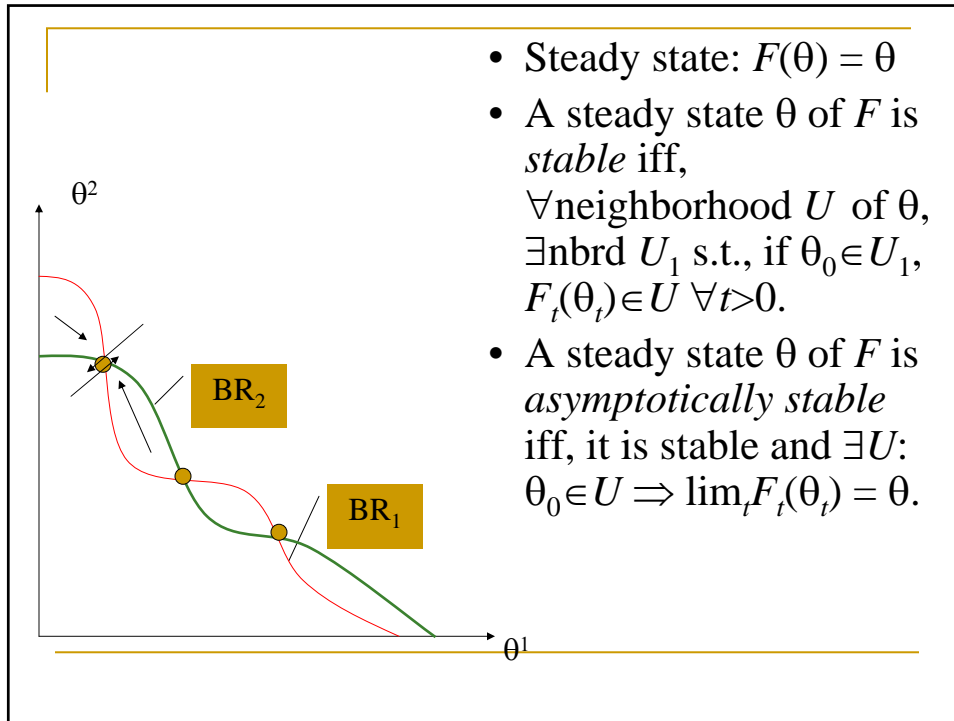
	R	S	P
R	0,0	1,-1	-1,1
S	-1,1	0,0	1,-1
P	1,-1	-1,1	0,0

- Unique Nash Equilibrium (s^*,s^*) where
 $s^* = (1/3,1/3,1/3)$
- s^* is not ESS.

ESS in role-playing games

- Given (S^1, S^2, u_1, u_2) , consider the symmetric game (\underline{S}, u) where
 - $\underline{S} = S^1 \times S^2$;
 - $u(\underline{x}, \underline{y}) = [u_1(x_1, y_2) + u_2(x_2, y_1)]/2 \quad \forall \underline{x} = (x_1, x_2), \underline{y} = (y_1, y_2) \in \underline{S}$.
- Theorem:** \underline{x} is an ESS of (\underline{S}, u) iff \underline{x} is a strict Nash equilibrium of (S^1, S^2, u_1, u_2) .

Replicator Dynamics



Replicator dynamics

- $p_i(t)$ = #people who plays s_i at t ;
 - $p(t)$ = total population at t .
 - $x_i(t) = p_i(t)/p(t)$; $x(t) = (x_1(t), \dots, x_k(t))$.
 - $u(x, x) = \sum_i x_i u(s_i, x)$.
 - Birthrate for s_i at t is $\beta + u(s_i, x(t))$; death rate = δ .
 - $\dot{p}_i = [\beta + u(s_i, x) - \delta] p_i$
 - $\dot{p} = [\beta + u(x, x) - \delta] p$
 - $\dot{x}_i = [u(s_i, x) - u(x, x)] x_i$
-
- $$\dot{x}_i = u(s_i - x, x) x_i$$

Example

- Consider (S,A) where $A = \begin{bmatrix} a_1 & 0 \\ 0 & a_2 \end{bmatrix}$
- $u(s_1, x) = a_1 x_1$;
- $u(x, x) = (x_1, x_2)A(x_1, x_2)^T = a_1 x_1^2 + a_2 x_2^2$
- $u(s_1 - x, x) = (a_1 x_1 - a_2 x_2)x_2$
- $\dot{x}_1 = (a_1 x_1 - a_2 x_2)x_1 x_2$

Examples

- Replicator dynamics in prisoners' dilemma
- Replicator dynamics in chicken
- Replicator dynamics in the battle of the sexes.

Observations

$$\begin{aligned} \blacksquare \frac{d}{dt} \left[\frac{x_i}{x_j} \right] &= \frac{\dot{x}_i}{x_j} - \frac{x_i}{x_j} \frac{\dot{x}_j}{x_j} = [u(s_i, x) - u(x, x)] \frac{x_i}{x_j} - \frac{x_i}{x_j} [u(s_j, x) - u(x, x)] \frac{x_j}{x_j} \\ &= [u(s_i, x) - u(s_j, x)] \frac{x_i}{x_j} \end{aligned}$$

- If u becomes $\underline{u} = au + b$, then Replicator dynamics becomes

$$\dot{x}_i = \underline{u}(s_i - x, x)x_i = au(s_i - x, x)x_i$$

Rationalizability

- $\xi(\cdot, x_0)$ is the solution to replicator dynamics starting at x_0 .

Theorem: If a pure strategy i is strictly dominated (by y), then $\lim_t \xi_i(t, x_0) = 0$ for any interior x_0 .

Proof: Define $v_i(x) = \log(x_i) - \sum_j y_j \log(x_j)$. Then,

$$\frac{dv_i(x(t))}{dt} = \frac{\dot{x}_i}{x_i} - \sum_j y_j \frac{\dot{x}_j}{x_j} = u(s_i - x, x) - \sum_j y_j u(s_j - x, x) = u(s_i - y, x).$$

Hence, $v_i(x(t)) \rightarrow -\infty$, i.e., $x_i(t) \rightarrow 0$.

Theorem: If i is not rationalizable, then $\lim_t \xi_i(t, x_0) = 0$ for any interior x_0 .

Theorems

Theorem: Every ESS x is an asymptotically stable steady state of replicator dynamics.

(If the individuals can inherit the mixed strategies, the converse is also true.)

Theorem: If x is an asymptotically stable steady state of replicator dynamics, then (x,x) is a perfect Nash equilibrium.

MIT OpenCourseWare
<http://ocw.mit.edu>

14.126 Game Theory
Spring 2010

For information about citing these materials or our Terms of Use, visit: <http://ocw.mit.edu/terms>.